

# On the Evaluation and Reduction of Generalized Polylogarithms

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## On the reduction of generalized polylogarithms to $\text{Li}_n$ and $\text{Li}_{2,2}$ and on the evaluation thereof

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**ABSTRACT:** We give expressions for all generalized polylogarithms up to weight four in terms of the functions  $\log$ ,  $\text{Li}_n$ , and  $\text{Li}_{2,2}$ , valid for arbitrary complex variables. Furthermore we provide algorithms for manipulation and numerical evaluation of  $\text{Li}_n$  and  $\text{Li}_{2,2}$ , and add codes in Mathematica and C++ implementing the results. With these results we calculate a number of previously unknown integrals, which we add in App. C.

**KEYWORDS:** Generalized polylogarithms, Multiple polylogarithms, Higher orders, Feynman integrals, Computer algebra.



## Introduction

Generalized Polylogarithms (GPLs) are a class of mathematical functions defined recursively as

$$G(a_1, \dots, a_n; x) = \int_0^x \frac{dt}{t - a_1} G(a_2, \dots, a_n; t)$$

with  $G(; x) \equiv 1$ .



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with  $G(; x) \equiv 1$ .

$$G(\underbrace{0, \dots, 0}_n; x) \equiv G(\bar{0}_n; x) = \frac{\log^n(x)}{n!}$$

$$G(0; x) = \log(x) \quad G(a; x) = \log\left(1 - \frac{x}{a}\right) \quad G(\bar{0}_{n-1}, a; x) = -\text{Li}_n\left(\frac{x}{a}\right)$$

The shuffle rule:

$$G(\bar{a}; x)G(\bar{b}; x) = \sum_{\bar{c} \in \bar{a}\amalg\bar{b}} G(\bar{c}; x)$$



## Motivation

$$G(a_1, \dots, a_n; x) = \int_0^x \frac{dt}{t - a_1} G(a_2, \dots, a_n; t) \quad G(0; x) = \log(x)$$
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For loop-corrections to scattering amplitude calculations, a lot of Feynman integrals are needed.

For each process, they can (using IBPs) be expressed in terms of a minimal set, denoted master integrals.



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Example:

The one-loop corection to massless  $2 \rightarrow 2$  scattering (e.g. QCD contribution to 2-jet production at the LHC)

Three master integrals: two 'bubbles', one 'box'.



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We define the vector  $f = \{I_{s\text{-bubble}}, I_{t\text{-bubble}}, I_{\text{box}}\}$  and  $x = s/t$ .  
Then  $\frac{df}{dx} = A(x, \epsilon)f$ .





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A smart choice of  $f$  gives the 'canonical' form (Henn: [arXiv:1304.1806])

$$df = \epsilon \sum_i d \log(q_i(x)) \tilde{A}_i f$$



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If we expand  $f(\epsilon, x) = \sum \epsilon^i f^{(i)}(x)$ , and if  $q_1 = x - a_1$ , then

$$\frac{df^{(n)}}{dx} = \frac{1}{x - a_1} \tilde{A}_1 f^{(n-1)} + \dots$$



## Reduction

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It was conjectured by Goncharov [arXiv:1110.0458] that all GPLs up to weight 4 (all that is needed at two-loop) can be expressed in terms of

$$\log(x), \text{Li}_2(x), \text{Li}_3(x), \text{Li}_4(x), \text{Li}_{2,2}(x, y)$$

$$\text{Li}_n(x) = \int_0^x \frac{dt}{t} \text{Li}_{n-1}(t) \quad \text{with} \quad \text{Li}_1(x) \equiv -\log(1-x)$$

$$\text{Li}_{2,2}(x, y) = \int_0^1 \frac{\log(t) \text{Li}_2(xyt)}{t - \frac{1}{x}} dt$$



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Let's show that...



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Reduction at weight two:  $A \neq B \neq X \neq 0$

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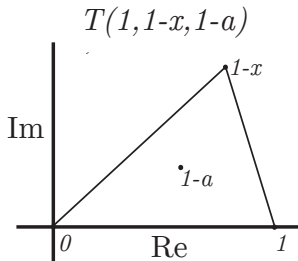
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Reduction at weight two (cont.):

$$\begin{aligned} G(A, B, X) &= G\left(\frac{A}{B}, 1; \frac{X}{B}\right) \equiv G(a, 1, x) \\ &= G(1 - a, 0; 1 - x) - G(1 - a, 0; 1) + 2\pi iT(1, 1 - x, 1 - a)G(0; 1 - a) \end{aligned}$$



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Goal achieved:  $G(a, b; x)$  in terms of  $\log$ ,  $\text{Li}_2$ .



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The reduction consists of two parts:

First map a letter to zero,  
then do the remaining integrals.

$$G(A_1, \dots, A_n; X) \rightarrow G(a_1, \dots, a_{n-1}, 0, 1)$$
$$\rightarrow \log, \text{Li}_n, \text{Li}_{2,2} \text{ at weights } \leq 4.$$

The first step can be done at all weights,  
The second gets much harder for each weight.



# Reduction

$$\begin{aligned}
 G(a, b, c, x) = & \frac{1}{6} \log \left( \frac{c-b}{a-x} \right)^3 - \frac{1}{6} \log \left( \frac{(a-c)(c-b)}{(a-b)c} \right)^3 - \frac{1}{6} \log \left( \frac{c-b}{a-x} \right)^3 + \frac{1}{6} \log \left( \frac{(a-c)(c-b)}{(a-b)(c-x)} \right)^3 - \frac{1}{2} \log \left( \frac{a}{a-b} \right) \log \left( \frac{b}{b-c} \right)^2 + \text{Li}_2 \left( \frac{x-b}{a-b} \right) \log \left( \frac{b-x}{b-c} \right) \\
 & - \frac{1}{2} \log \left( 1 - \frac{b}{a} \right) \log \left( \frac{b}{b-c} \right)^2 + \frac{1}{2} \log \left( \frac{a-x}{a-b} \right) \log \left( \frac{b-x}{b-c} \right)^2 + \frac{1}{2} \log \left( \frac{a-x}{a-b} \right) \log \left( \frac{b-x}{b-c} \right) + \text{Li}_3 \left( \frac{b}{a} \right) + \text{Li}_3 \left( \frac{b}{a-b} \right) - \text{Li}_3 \left( \frac{b(a-c)}{a(b-c)} \right) - \text{Li}_3 \left( \frac{b(a-c)}{(a-b)c} \right) \\
 & - \text{Li}_3 \left( -\frac{c}{b-c} \right) + \text{Li}_3 \left( 1 - \frac{c}{a} \right) - \text{Li}_3 \left( \frac{a-c}{a-x} \right) - \text{Li}_3 \left( \frac{b-x}{a-x} \right) + \text{Li}_3 \left( \frac{(a-c)(b-x)}{(b-c)(a-x)} \right) + \text{Li}_3 \left( \frac{(a-c)(b-x)}{(a-b)(c-x)} \right) - \text{Li}_3 \left( \frac{x-b}{a-b} \right) - \text{Li}_3 \left( \frac{x-c}{a-c} \right) + \text{Li}_3 \left( \frac{x-c}{b-c} \right) \\
 & - \text{Li}_2 \left( \frac{a-c}{b-c} \right) \log \left( \frac{a}{a-c} \right) + \text{Li}_2 \left( \frac{-c}{b-c} \right) \log \left( \frac{a}{a-c} \right) - \text{Li}_2 \left( \frac{b}{a} \right) \log \left( \frac{b}{b-c} \right) - \text{Li}_2 \left( \frac{b}{a-b} \right) \log \left( \frac{b}{b-c} \right) + \text{Li}_2 \left( \frac{b(a-c)}{a(b-c)} \right) \log \left( \frac{b}{b-c} \right) + \text{Li}_2 \left( \frac{b(a-c)}{(a-b)c} \right) \log \left( \frac{b}{b-c} \right) \\
 & + \frac{1}{2} \log \left( \frac{b}{b-c} \right)^2 \log \left( \frac{(b-a)c}{a(b-c)} \right) + \frac{1}{2} \log \left( \frac{b}{b-c} \right)^2 \log \left( \frac{a(c-b)}{(a-b)c} \right) + \text{Li}_2 \left( \frac{a-c}{b-c} \right) \log \left( \frac{a-x}{a-c} \right) - \text{Li}_2 \left( -\frac{c}{b-c} \right) \log \left( \frac{a-x}{a-c} \right) + \text{Li}_2 \left( \frac{b-x}{a-x} \right) \log \left( \frac{b-x}{b-c} \right) \\
 & - \text{Li}_2 \left( \frac{(a-c)(b-x)}{(b-c)(a-x)} \right) \log \left( \frac{b-x}{b-c} \right) - \text{Li}_2 \left( \frac{(a-c)(b-x)}{(a-b)(c-x)} \right) \log \left( \frac{b-x}{b-c} \right) - \frac{1}{2} \log \left( \frac{b-x}{b-c} \right)^2 \log \left( -\frac{(b-c)(a-x)}{(a-b)(c-x)} \right) - \frac{1}{2} \log \left( \frac{b-x}{b-c} \right)^2 \log \left( \frac{(b-c)(c-x)}{(b-c)(a-x)} \right) + \text{Li}_3 \left( \frac{c}{c-a} \right) \\
 & - \text{Li}_2 \left( \frac{c-b}{a-b} \right) \log \left( 1 - \frac{x}{c} \right) + \text{Li}_2 \left( \frac{x-b}{a-b} \right) \log \left( 1 - \frac{x}{c} \right) + \log \left( \frac{a-x}{a-b} \right) \log \left( \frac{b-x}{b-c} \right) \log \left( 1 - \frac{x}{c} \right) - \pi^2 \left( -\frac{1}{6} \log \left( \frac{c-b}{a-b} \right) + \frac{1}{6} \log \left( \frac{(a-c)(c-b)}{(a-b)c} \right) + \frac{1}{6} \log \left( \frac{c-b}{a-x} \right) \right. \\
 & \left. - \frac{1}{6} \log \left( \frac{(a-c)(c-b)}{(a-b)(c-x)} \right) + 4 \log \left( 1 - \frac{b}{c} \right) \text{T} \left( 1, 1 - \frac{x}{c}, 1 - \frac{b}{c} \right) \text{T} \left( \text{P} \left( \frac{b}{c}, 1 - \frac{x}{c} \right), 1 - \frac{x}{c}, 1 - \frac{a}{c} \right) + i\pi \left( \text{T} \left( 1, \frac{b}{c}, \frac{a}{a-c} \right) \text{sgn} \left( \text{Im} \left( \frac{c}{a-c} \right) \right) \log \left( \frac{a}{a-c} \right)^2 \right. \right. \\
 & \left. \left. + \log \left( \frac{b-a}{b-c} \right)^2 \text{T} \left( 1, \frac{b}{c}, \frac{b-a}{b-c} \right) \text{sgn} \left( \text{Im} \left( \frac{a-c}{b-c} \right) \right) - \log \left( \frac{b-a}{b-c} \right)^2 \text{T} \left( 1, \frac{b-x}{b-c}, \frac{b-a}{b-c} \right) \text{sgn} \left( \text{Im} \left( \frac{a-c}{b-c} \right) \right) - 2\text{Li}_2 \left( \frac{a-c}{b-c} \right) \text{T} \left( 1, 1 - \frac{x}{c}, 1 - \frac{a}{c} \right) \text{sgn} \left( \text{Im} \left( \frac{a}{c} \right) \right) \right. \right. \\
 & \left. \left. + 2\text{Li}_2 \left( -\frac{c}{b-c} \right) \text{T} \left( 1, 1 - \frac{x}{c}, 1 - \frac{a}{c} \right) \text{sgn} \left( \text{Im} \left( \frac{a}{c} \right) \right) - 2 \log \left( 1 - \frac{a}{c} \right) \log \left( \frac{b-a}{b-c} \right) \text{T} \left( 1, 1 - \frac{x}{c}, 1 - \frac{a}{c} \right) \text{sgn} \left( \text{Im} \left( \frac{a}{c} \right) \right) \right. \right. \\
 & \left. \left. - \mathcal{H}_1 \left( 1 - \frac{a}{c}, 1 - \frac{b}{c} \right) \log \left( \frac{(b-a)c}{(a-c)(c-b)} \right)^2 \text{sgn} \left( \text{Im} \left( \frac{b}{c} \right) \right) + 2 \log \left( 1 - \frac{b}{c} \right) \log \left( \frac{a-b}{a-c} \right) \text{T} \left( 1, 1 - \frac{x}{c}, 1 - \frac{b}{c} \right) \text{sgn} \left( \text{Im} \left( \frac{b}{c} \right) \right) \right. \right. \\
 & \left. \left. - 2 \log \left( 1 - \frac{b}{c} \right) \log \left( \frac{a-x}{a-c} \right) \text{T} \left( 1, 1 - \frac{x}{c}, 1 - \frac{b}{c} \right) \text{sgn} \left( \text{Im} \left( \frac{b}{c} \right) \right) + \mathcal{H}_1 \left( \frac{b-c}{a-c}, \frac{c-b}{c-x} \right) \log \left( \frac{a-x}{b-c} \right)^2 \text{sgn} \left( \text{Im} \left( \frac{c-b}{c-x} \right) \right) \right. \right. \\
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$\text{Li}_n(x)$  is the classical polylogarithm (or Euler polylogarithm)

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Outside we may use the inversion relation  $\text{Li}_n(x) = \pm \text{Li}_n(\frac{1}{x}) + \text{easier}$ .

$$\text{Li}_3(x) = \text{Li}_3\left(\frac{1}{x}\right) - \frac{1}{6} \log^3(-x) - \frac{\pi^2}{6} \log(-x)$$

(There are better ways to evaluate  $\text{Li}_n(x)$ ,  
e.g. a 'Bernoulli expansion'  $\text{Li}_n(e^x) = \sum_{\nu} \frac{c_{\nu}}{\nu!} x^{\nu}$ ).



$\text{Li}_{2,2}(x, y)$  is given by a similarly looking sum

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If  $|xy| > 1$  we may use the inversion relation

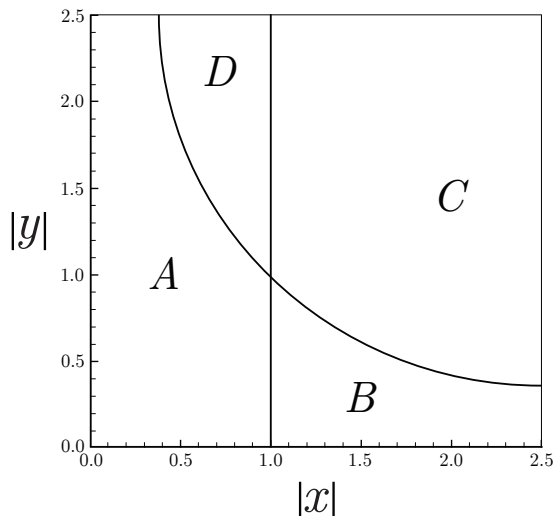
$$\begin{aligned} \text{Li}_{2,2}(x, y) &= \text{Li}_{2,2}\left(\frac{1}{x}, \frac{1}{y}\right) - \text{Li}_4(xy) + 3\left(\text{Li}_4\left(\frac{1}{x}\right) + \text{Li}_4(y)\right) \\ &+ 2\left(\text{Li}_3\left(\frac{1}{x}\right) - \text{Li}_3(y)\right) \log(-xy) + \text{Li}_2\left(\frac{1}{x}\right) \left(\frac{\pi^2}{6} + \frac{1}{2} \log^2(-xy)\right) \\ &+ \frac{1}{2} \text{Li}_2(y) \left(\log^2(-xy) - \log^2(-x)\right) \end{aligned}$$

And then if  $|x| > 1$  we may use the 'stuffle relation'

$$\text{Li}_{2,2}(x, y) = -\text{Li}_{2,2}(y, x) - \text{Li}_4(xy) + \text{Li}_2(x)\text{Li}_2(y)$$



## Evaluation



*A:* Direct summation

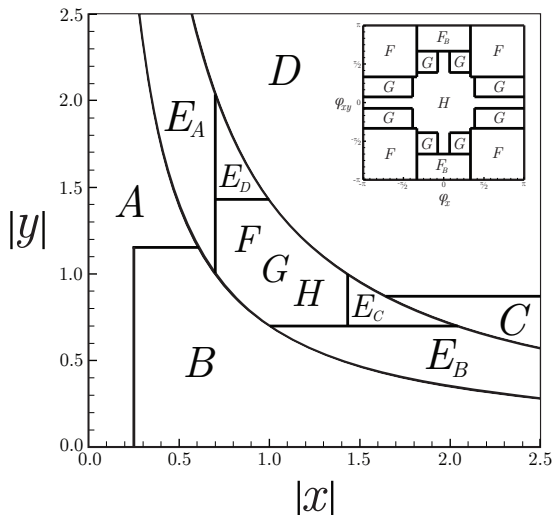
*B:* Stuffle:  $x \leftrightarrow y$

*C:* Inversion:  $z \leftrightarrow z^{-1}$

*D:* Stuffle and inversion



# Evaluation



- A: Direct summation
- B: Stuffle:  $x \leftrightarrow y$
- C: Inversion:  $z \leftrightarrow z^{-1}$
- D: Stuffle and inversion
- E: Diagonal sum
- F: Hölder relation
- G:  $\log(1 - z)$  expansion
- H:  $\log(z)$  expansion



## The added code

With the paper [ArXiv:1601.02649] we have added code implementing the results.

The reductions are implemented in Mathematica as the replacement rule `gtolrules`.

```
In[1] := G(0,a,0,b,x)/.gtolrules  
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Please try our code, and send us any critique/bug reports.



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The past decade has seen many new developments in the study of GPLs.

Particularly the introduction of 'symbols' and co-products  
[arXiv:1006.5703, arXiv:1110.0438, arXiv:1123.0454]

$$\mathcal{S}\left(G(a_n, \dots, a_1; a_{n+1})\right) = \sum_{i=1}^n \left( \mathcal{S}\left(G(a_n, \dots, \hat{a}_i, \dots, a_1; a_{n+1})\right) \otimes (a_i - a_{i+1}) \right. \\ \left. - \mathcal{S}\left(G(a_n, \dots, \hat{a}_i, \dots, a_1; a_{n+1})\right) \otimes (a_i - a_{i-1}) \right)$$



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Symbols capture the algebraic parts of an expression.  
We need the full analytical structure.

Newer developments of the co-products catch most  
of the analyticity [arXiv:1401.3546]

To generalize to weight 6 (or  $n$ ) this will likely be needed.



The goal of reducing all GPLs up to weight four to

$$\log(x) , \operatorname{Li}_2(x) , \operatorname{Li}_3(x) , \operatorname{Li}_4(x) , \operatorname{Li}_{2,2}(x, y)$$

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Additionally a robust implementation of  $\operatorname{Li}_{2,2}(x, y)$  was made.



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Thank you for listening...

