

# Storage

In HEP,

on the Grid,

and in the Computing Center

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GridKa-School, HEP session

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# HEP and storage use

## > HEP has always been major storage user

- Bubble chambers ... (HERA, BaBar, Belle, Tevatron, RHIC...)

## > With LHC: HEP is going into a new dimension

- LHC has pushed Data Grid technologies
- Vendors embracing HEP datacenters, good clients:-)

## > Future: Not clear whether HEP still has a key position

- Data mining, data warehousing, ... rapidly growing capacity needs in industry
- Other science (e.g. XFEL@DESY) similar data rates expected than HEP
- **First lesson:** HEP should stick to industry standards!



# The LHC data challenge and the LHC Grid

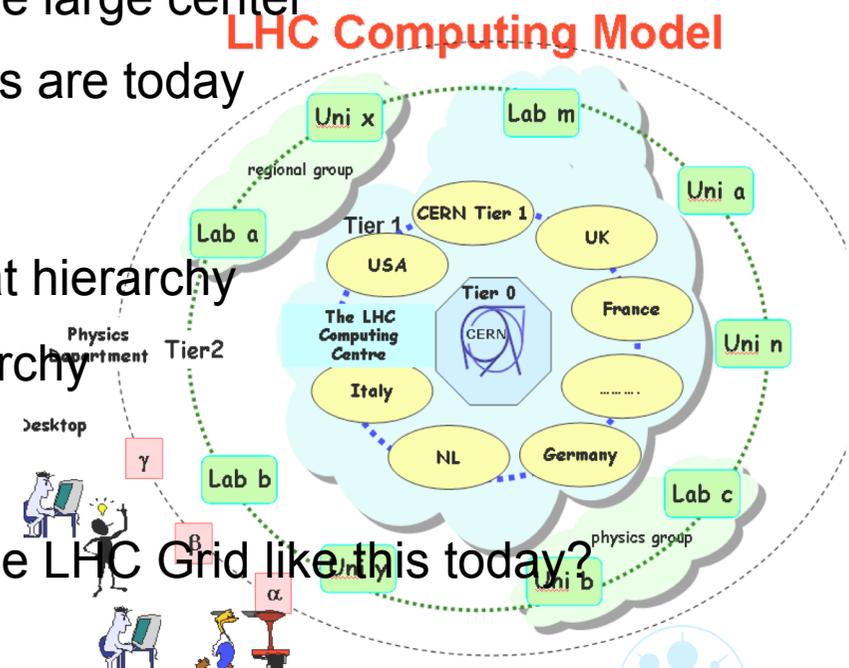
- > (you know the numbers, no need to bring them up again)
- > Hosting CPUs and storage in one single/very few places impossible
  - Technology: Dimension just far beyond current state of art
  - Network: Is slow and expensive
  - Security: Need at least a second place and copy in case the first one breaks
  - Politics: Did not want to put all money into one large center
  - That was back in 1990ths... guess how things are today

## > Two things come together ...

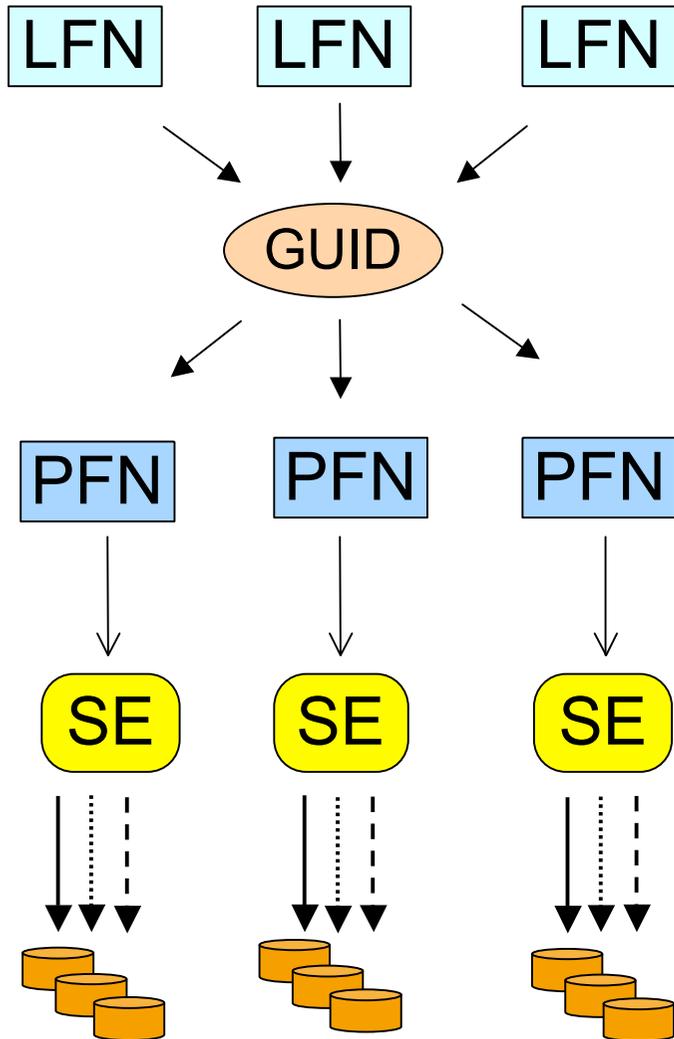
- The Computing Grid with its (theoretically) flat hierarchy
- A tiered layer of data centers with clear hierarchy

## > ... to create the LHC Grid

- An heretical question: Would one still build the LHC Grid like this today?



# Coming from the Grid view to the local view



## > “Filename” on the Grid

- Global Unique Identifier (GUID)
- `guid:3a69a819-2023-4400-a2a1-f581ab942044`

## > Easier with Logical File Name (LFN)

- `lfn:/grid/myexp/kemp/ExitingDataset.dat`
- `lfn:/grid/myexp/myboss/DataWithBadDetector.dat`

## > Physical File Name (PFN)

- A Path on the SE: also called Storage URL (SURL)
- `/storage/grid/experiments/cms/kemp/ver04/run2342/results/data/file124.dat`
- Files can be replicated to several SEs

## > Up to here: Correspondence governed by catalogues

- Like LFC

## > From the PFN to the transport URL (TURL)

- The transport protocol: `(gsi)dcap`, `gsifp`, `xrootd`
- The SE (SRM) will tell you
- ... and you will access the data on the hardware....

# The “Ideal” storage: Attributes to storage

## > Fast:

- Fast in getting the meta-data (“ls -l”)
- Fast in getting the first desired bit / random reads
- Fast in getting a sustained stream
- Fast in writing the data

## > Huge

## > Unique, consistent and easy:

- Unique namespace (no “/tmp”, “/afs/...”, “/grid/cms/...”, “/home”, “srm://pnfs...”)
- Consistent access methods throughout the whole storage
- Easy access to the data

## > Accessible from everywhere

- And fast ;-)

## > Cheap

- Purchase
- Running costs
- cooling, electricity, space consumption, ...



Harry Potter tm Trunk with Dressing Up Set.

# Some more attributes...

- > Secure (data integrity)
  - Authentication and authorization (no one else can temper your data)
  - Backup (even you cannot temper your data by mistake)
  - Robust media/technology & backup (even a disaster cannot temper your data)
- > Simple manageability, stable running, good support
  - Little administration costs, good vendor support
  - Little disturbances by downtimes / maintenance
- > Migration
  - If a newer / better system becomes available: No vendor lock-in
- > Long term availability
  - Of your data
  - Of the storage system
  - Of the protocols
- > .....

## Conclusion:

*“The One Ideal”* storage does not exist  
Compromises, and different products  
for different purposes



# Some technology: Media

## > RAM and NVRAM (e.g. battery powered DRAM (+disks))

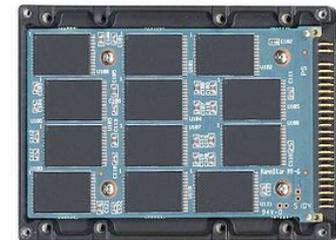
- Yes: RAM-Disks do exist: Databases!
- Sometimes used as Meta-Data disks for file servers in HEP

## > Solid State Disks

- Have emerged in the last year, become less and less expensive
- Serious competitor to Hard-Drives in some future
- Different access behavior than traditional “spindle disks”

## > Hard Disk Drives (with magnetic spindles)

- Established technology
- High density, and increasing
- Streaming performance very good
- Random access / seek time relatively slow w.r.t. streaming
- Different connections / qualities: P-ATA, S-ATA, SCSI, SAS, FC, ...



Wikimedia Commons



# More technology



## > Tapes

- “Will disappear soon”: Sentence true since (at least) 10 years :-)
- And still tape is the working horse for storing data at CERN, FZK, DESY and elsewhere
- Lowest media cost (~50 EUR / TB), Green-IT (no electricity when not accessed)
- Best scaling storage system available, difficult to handle (administration, access, ....)

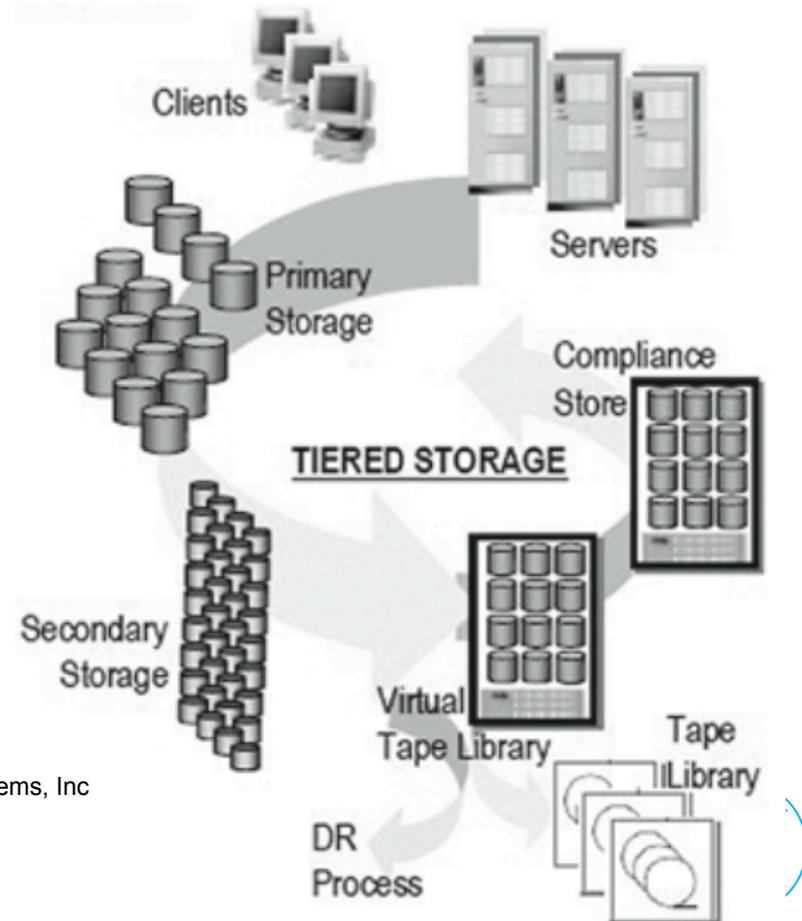
## > Optical media (CD, DVD, ...)

- Play only a minor role in large scale data storage

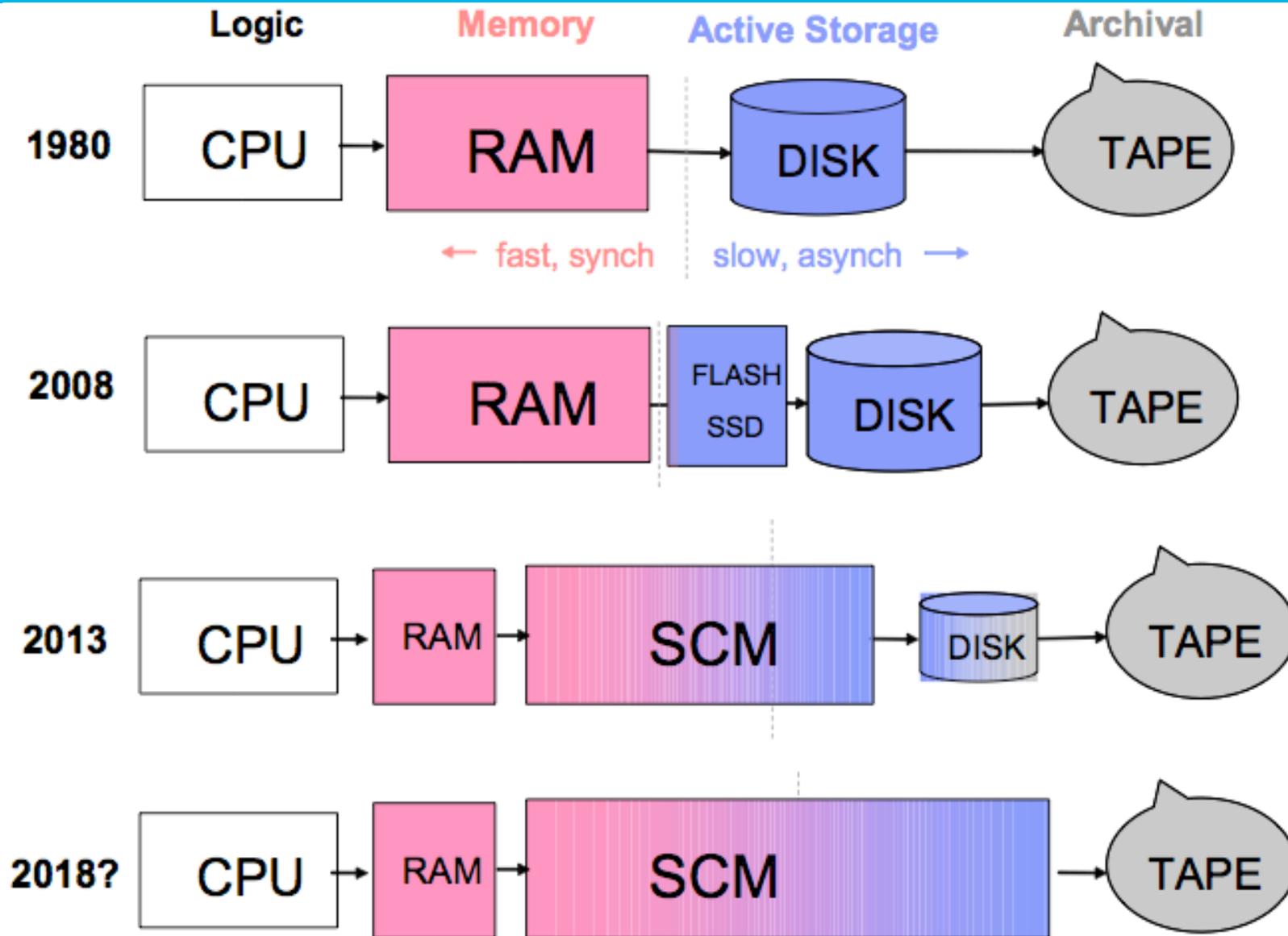
Putting them together:

## > Tiered storage

- Migration from active disks to offline storage
- Automatically, transparent to users



# Storage Class Memory



Slide: David A. Pease, IBM



# One interesting thing about RAID and data security

- > RAID: Redundant Array of Independent|Inexpensive Disks
  - RAID-0: Stripe information among >1 disks
  - RAID-1: Mirror information on 2 disks
  - RAID-5: Capacity of N disks, using N+1 disks, one disk used for checksums
  - RAID-6: Capacity of N disks, using N+2 disks, two disks used for checksums
  - RAID-10/50... combinations of the above
- > HDD error rate between  $10^{-16}$  and  $10^{-14}$ 
  - Results in reading error every 10-1000 TB
- > RAID-5 is secure ... unless you store more than ~10 TB:-)
  - RAID-6 is somewhat more secure
- > **Absolute security does not exist! Some data will get lost!**
  - Also tapes (l.e. backup) can fail!



# Different storage places in the NAF

## > OK, what is the optimal workplace on the NAF?

- No single answer, but you get my personal recommendations for free :-)

## > AFS /afs/naf.desy.de/...: Network file system

- Login files, small data amounts (like plot.root) (total <1 GB)
- Source files for code (exclude libs or bins: Check HowTo with your VO admins)
- There are group volumes for SW releases. Check with your VO admins
- Compilation can be slow (Atlas-CMT problem), usually OK
- Not available on Grid-WNs (not directly, and I will not tell you how)
- Backup

## > /tmp: Local file system

- Is local, quite large, no quota, somewhat fast
- Cleaned up every 10 days

## > /scratch/... (Currently Lustre file system)

- Fast cluster file system, available everywhere in NAF but not externally, no backup, (currently) no quota or ACLs
- Using InfiniBand as interconnect, low latency and high bandwidth
- Currently optimized for large files, bad for unpacking source code and compilation
- Useful for temporary storage of “hot data”
- Storage of often-used personal NTuples

## > dCache / SE

- Central (Grid) Import/Export system, well integrated into experiment’s workflows
- Large data sets, shared by many people and accessible from everywhere
- Not “filesystem-like”: No compilation etc.
- Backup / Archive possible (not done yet)

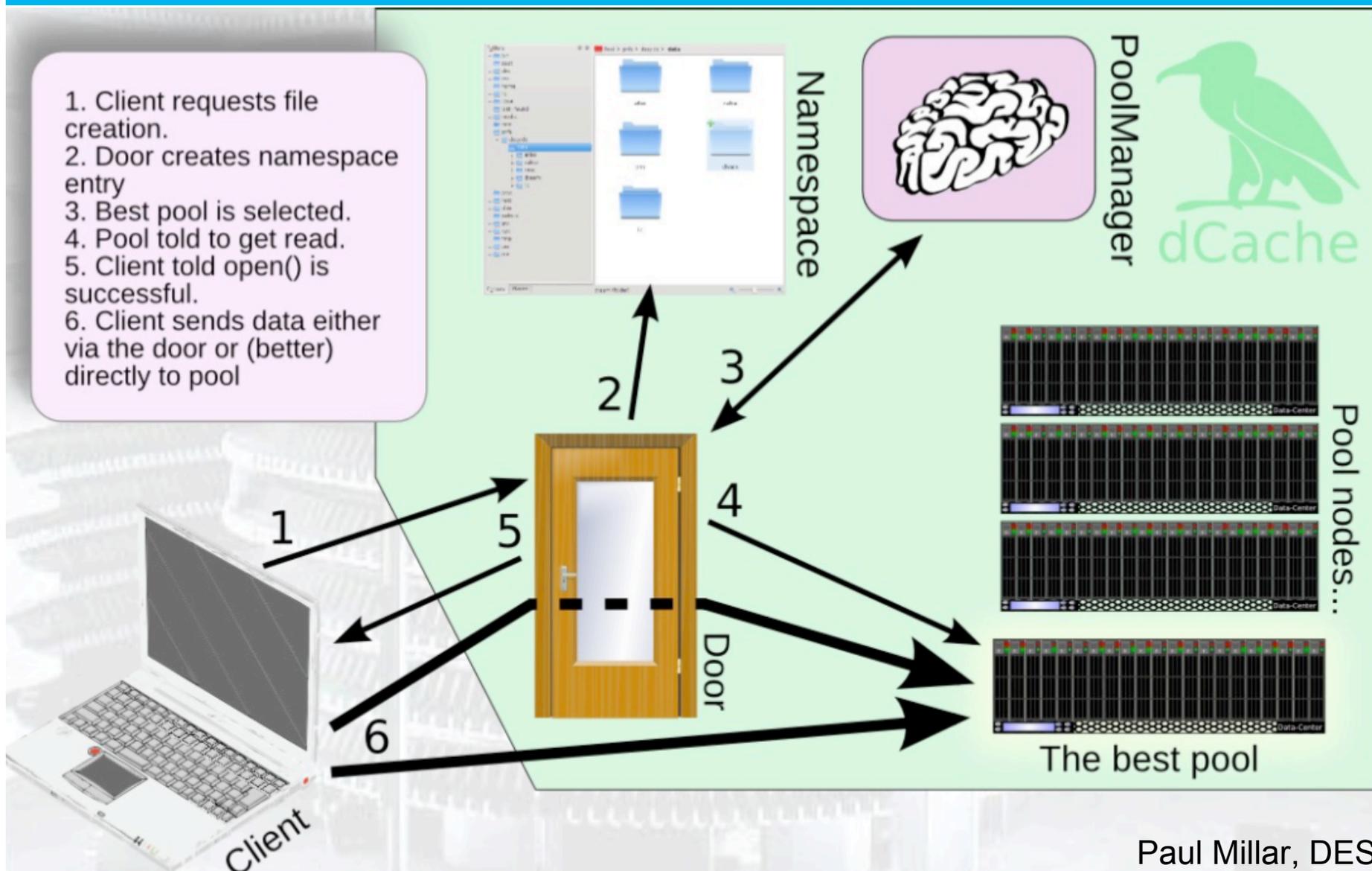


# dCache in a nutshell

- > dCache was introduced as a disk cache for tapes
  - Combine large but slow tape systems with small but fast file servers
  - Remember the “tiered-storage” plot???
- > Today can manage up to 10 PB of data
  - But also “Tier-3-like” installations, with  $O(100 \text{ TB})$  of data
- > Speaks many languages
  - (grid-)ftp
  - dcap, gsidcap
  - Xrootd protocol
  - http
  - SRM as a meta-language
- > Other systems (CASTOR, DPM, ...) have similar setup



# Workflow : writing files

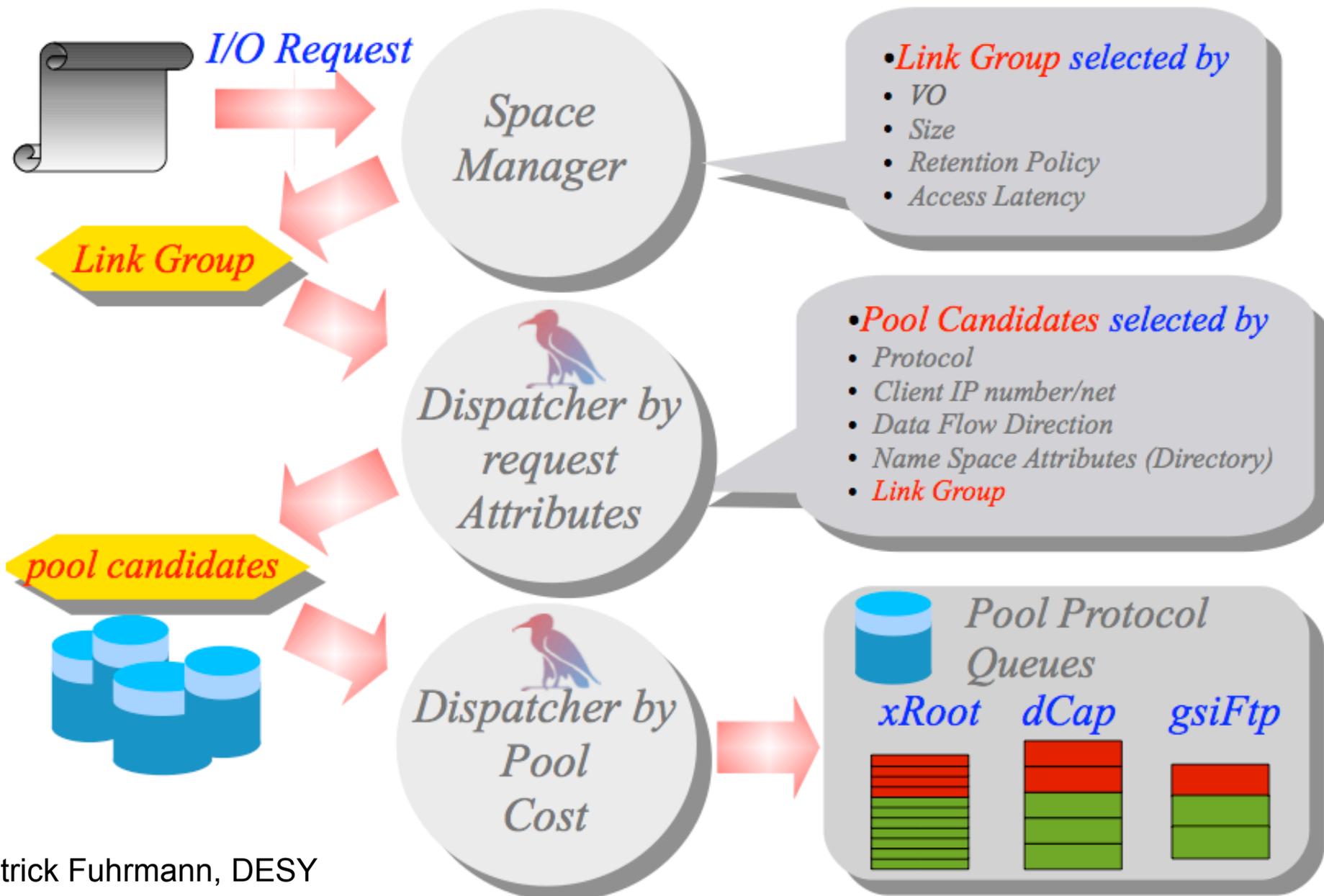


Paul Millar, DESY



# dCache a coarse view inside

## Request Flow



# SRM: Storage Resource Management

> Network protocol providing abstraction layer

> What SRM does:

- Negotiates transfer proto
- File pinning / unpinning
- Space management
- Name-space operations
- Permission management

> What SRM does not do:

- Data transfer
- Configuring data placement / policy engines
- Provisioning

> Two file attributes:

- Access Latency is: online, nearline, (offline)
- Retention Policy is: replica, (output), custodial.



<http://iris-ict.eu/joomla/images/stories/storage7.jpg>

# dCache: Pools and Doors

## > Pools: Hold the data

- Poolgroup: Group of pools:-) (e.g. all pools from mcdisk, data09disk,...)
- Poolnode: A computer with pools on it (usually fileserver, or attached to SAN)
- Replicates: One file can have several copies on different pools (of the same poolgroup). Can be done automatically, useful for increasing performance

## > Doors: Connect you with your data

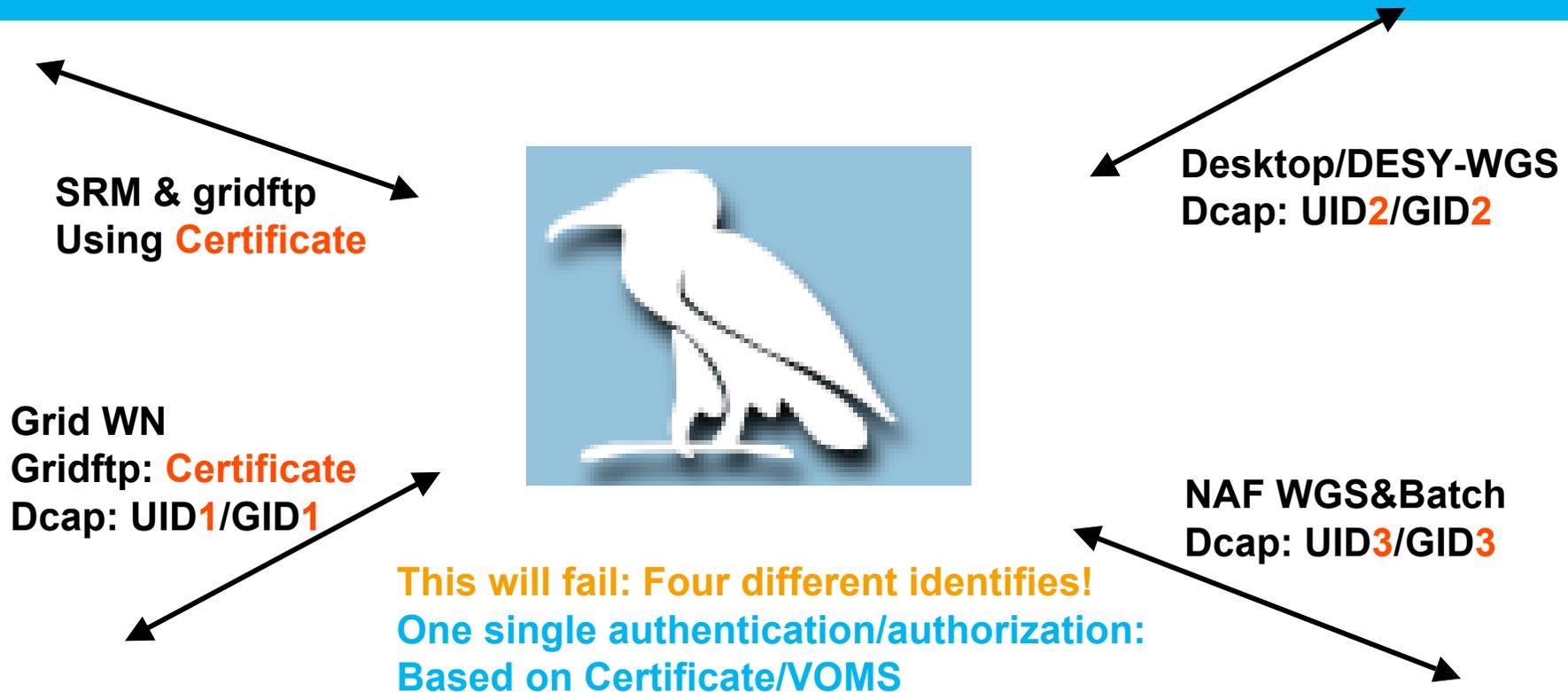
- Several doors, for each protocol: (gsi)dcap, (gsi)ftp, http(s), xrootd, ...
- Can have more than one / protocol
- Ask SRM for the door with protocol X → You will get the best matching door
- You communicate with this best matching door ...

## > Getting/Putting the data

- Either your client is redirected to the pool, and gets/puts data directly
- ((Or the access goes via the door, good for firewalls, bad for speed))



# Authentication/Authorization



## > Protocols

- Gridftp (same as before)
- Gsidcap: Same as dcap, but with GSI authz

## > E.g. ROOT supports gsidcap

## > Meta-Data handling (e.g. file browsing)

- /pnfs/ needs dcap (not gsidcap!)
- dcTools developed at DESY by summer student Malte Nuhn
- On NAF: ini dctools -> dcls -l /pnfs/desy.de/ilc

## > SRM tools also OK, but slower

## > http based solution under development

# Speed discussion

- > Hepix storage task force: dCache similar speed than other products
  - DPM, XROOTD server and dCache: No seizable difference in performance
  - Will take dCache as an example in the following
- > dCache has movers, and they might get queued if too many
  - To protect the system against overload
  - You might have to wait :- ) (We see this, and can optimize things up to some point)
- > SRM is slow. People know this. Unfortunately, there is not much we can do about it...
- > Communication overhead (doors, pools,...)
  - ( $\sim 0.5$  s communication / file open) + ( $\sim 0.5$  s GSI security / session)
  - dCache developers try to lower both overheads
- > Data transfer is very fast
  - In streaming mode, near to wire speed ((gsi)dcap and gsiftp)



# How fast a system must be?

## > Example math:

- One job: 10 events/s
- One event (AOD): 150 kByte
- → 1,5 MByte/s/job (CPU limited....)
- 5 million events per dataset
- Want to compute this in 1 hour

## > $5 \cdot 10^7 / 10 / 3600 = 140$ jobs in parallel

## > 140 jobs \* 1,5 MByte = 210 MByte/s aggregate bandwidth

- We have tested Lustre with 500 MByte/s (one server!)
- dCache: Data distributed over different pools. Now: ~100 MByte/s/poolnode, soon: up to 1 GByte/s/poolnode

## > Numbers can/will/should change!



M. Schumacher on Ferrari, 2005  
Wikimedia commons



# The Last Slide

**No conclusion here** (except that this is an incomplete talk with a lot of personal opinions)

## Any questions? Suggestions?

One question to you: We are always looking for benchmark applications. If you got a physics analysis application and are willing to spend some time rerunning your app against several storage technologies and different configurations, please contact us!

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