Programming Concepts in Mathematica

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Computer Algebra Systems

- Commercial systems: Mathematica, Maple, Matlab, MuPAD/Matlab, MathCad, Fermat, Derive...
- Free systems: FORM, GiNaC, Axiom, Cadabra, GAP, Reduce, Singular, Maxima, MAGMA...
- Generic systems: Mathematica, Maple, MuPAD/Matlab, Maxima, MathCad, Reduce, Axiom, MAGMA, GiNaC...
- Specialized systems: Cadabra, Singular, Magma, CoCoA, GAP...
- Many more...

Expert Systems

In technical terms, Mathematica is an Expert System. Knowledge is added in form of Transformation Rules. An expression is transformed until no more rules apply.

Example:

myAbs[x_] := x /; NonNegative[x]
myAbs[x_] := -x /; Negative[x]

We get:

Immediate and Delayed Assignment

Transformations can either be

• added "permanently" in form of Definitions,

• applied once using Rules:

a + b + c /. a -> 2 c 🖙 b + 3 c

Transformations can be Immediate or Delayed. Consider:

Mathematica is one of those programs, like TEX, where you wish you'd gotten a US keyboard for all those braces and brackets.

Almost everything is a List

All Mathematica objects are either Atomic, e.g.

Head[133] @ Integer

or (generalized) Lists with a Head and Elements:

expr = a + b
FullForm[expr] Plus[a, b]
Head[expr] Plus
expr[[0]] Plus - same as Head[expr]
expr[[1]] a
expr[[2]] b

List-oriented Programming

Using Mathematica's list-oriented commands is almost always of advantage in both speed and elegance.

Consider:

```
array = Table[Random[], {10^7}];
```

```
test1 := Block[ {sum = 0},
   Do[ sum += array[[i]], {i, Length[array]} ];
   sum ]
```

```
test2 := Apply[Plus, array]
```

Here are the timings:

Timing[test1][[1]] ③ 31.63 Second Timing[test2][[1]] ③ 3.04 Second Map, Apply, and Pure Functions

Map applies a function to all elements of a list:
 Map[f, {a, b, c}] @ {f[a], f[b], f[c]}
 f /0 {a, b, c} @ {f[a], f[b], f[c]} - short form

Apply exchanges the head of a list: Apply[Plus, {a, b, c}] \Rightarrow a + b + c Plus @@ {a, b, c} \Rightarrow a + b + c — short form

Pure Functions are a concept from formal logic. A pure function is defined 'on the fly':

(# + 1)& /@ {4, 8} 🖙 {5, 9}

The # (same as #1) represents the first argument, and the & defines everything to its left as the pure function.

List Operations

Flatten removes all sub-lists: Sort and Union sort a list. Union also removes duplicates: **Prepend and Append add elements at the front or back:** Prepend[r[a, b], c] r[c, a, b]Append[r[a, b], c] r[a, b, c]**Insert** and **Delete** insert and delete elements: Insert[h[a, b, c], x, {2}] <>>> h[a, x, b, c]

Patterns

One of the most useful features is **Pattern Matching**:

matches one object
 matches one or more objects
 matches zero or more objects
 matches zero or more objects
 named pattern (for use on the r.h.s.)
 x_h
 pattern with head h
 x_:1
 default value
 x_?NumberQ
 conditional pattern
 x_ /; x > 0

Patterns take function overloading to the limit, i.e. functions behave differently depending on *details* of their arguments:

Attributes[Pair] = {Orderless}
Pair[p_Plus, j_] := Pair[#, j]& /@ p
Pair[n_?NumberQ i_, j_] := n Pair[i, j]

Attributes

Attributes characterize a function's behaviour before and while it is subjected to pattern matching. For example,

Attributes[f] = {Listable}
f[1_List] := g[1]
f[{1, 2}] f[1], f[2]} - definition is never seen

Important attributes: Flat, Orderless, Listable, HoldAll, HoldFirst, HoldRest.

The Hold... attributes are needed to pass variables by reference:

Attributes[listadd] = {HoldFirst}
listadd[x_, other_] := x = Flatten[{x, other}]

This would not work if x were expanded before invoking listadd, i.e. passed by value.

Memorizing Values

For longer computations, it may be desirable to 'remember' values once computed. For example:

```
fib[1] = fib[2] = 1
fib[i_] := fib[i] = fib[i - 2] + fib[i - 1]
fib[4]  3
fib[4]  3
?fib  Global'fib
   fib[1] = 1
    fib[2] = 1
    fib[2] = 1
   fib[3] = 2
   fib[4] = 3
   fib[4] = 3
   fib[i_] := fib[i] = fib[i - 2] + fib[i - 1]
```

Note that Mathematica places more specific definitions before more generic ones.

Decisions

Mathematica's **If Statement** has three entries: for True, for False, but also for Undecidable. For example:

If[8 > 9, yes, no] If[a > b, yes, no] If[a > b, yes, no] If[a > b, yes, no]
If[a > b, yes, no, dunno] If[a > b, yes, no, dunno]

Property-testing Functions end in Q: EvenQ, PrimeQ, NumberQ, MatchQ, OrderedQ, ... These functions have no undecided state: in case of doubt they return False.

Conditional Patterns are usually faster:

Equality

Just as with decisions, there are several types of equality, decidable and undecidable:

a === b ③ a === b a === b ③ False a === a ④ True a === a ☞ True

The full name of '===' is SameQ and works as the Q indicates: in case of doubt, it gives False. It tests for Structural Equality. Of course, equations to be solved are stated with '==': Solve [$x^2 == 1$, x] $\Rightarrow \{\{x \rightarrow -1\}, \{x \rightarrow 1\}\}$ Needless to add, '=' is a definition and quite different: x = 3 — assign 3 to x

Selecting Elements

Select selects elements fulfilling a criterium: **Cases** selects elements matching a pattern: Using Levels is generally a very fast way to extract parts: list = {f[x], 4, {g[y], h}} **Depth** [list] \Leftrightarrow 4 — list is 4 levels deep (0, 1, 2, 3) Level[list, $\{3\}$] \Leftrightarrow $\{y\}$ Cases[expr, _Symbol, {-1}]//Union - find all variables in expr T. Hahn, Programming Concepts in Mathematica - p.14

Mathematical Functions

Mathematica is equipped with a large set of mathematical functions, both for symbolic and numeric operations. Some examples:

Sum[i, {i,50}] Series [Sin[x], $\{x,1,5\}$] Simplify $[(x^2 - x y)/x]$ Together [1/x + 1/y]Inverse[mat] Eigenvalues[mat] PolyLog[2, 1/3]LegendreP[11, x] Gamma[.567]

- integral
- derivative
- sum
- series expansion
- simplify
- put on common denominator
- matrix inverse
- eigenvalues
- polylogarithm
- Legendre polynomial
- Gamma function

Graphics

Mathematica has formidable graphics capabilities:

Plot[ArcTan[x], {x, 0, 2.5}]
ParametricPlot[{Sin[x], 2 Cos[x]}, {x, 0, 2 Pi}]
Plot3D[1/(x^2 + y^2), {x, -1, 1}, {y, -1, 1}]
ContourPlot[x y, {x, 0, 10}, {y, 0, 10}]

Output can be saved to a file with Export:

plot = Plot[Abs[Zeta[1/2 + x I]], {x, 0, 50}]
Export["zeta.eps", plot, "EPS"]

Hint: To get a high-quality plot with proper LATEX labels, don't waste your time fiddling with the Plot options. Use the psfrag LATEX package.

Numerics

Mathematica can express Exact Numbers, e.g.

Sqrt[2], Pi, $\frac{27}{4}$

It can also do Arbitrary-precision Arithmetic, e.g.

N[Erf[28/33], 25] \Rightarrow 0.7698368826185349656257148

But: Exact or arbitrary-precision arithmetic is fairly slow! Mathematica uses Machine-precision Reals for fast arithmetic.

Arrays of machine-precision reals are internally stored as Packed Arrays (this is invisible to the user) and in this form attain speeds close to compiled languages on certain operations, e.g. eigenvalues of a large matrix.

Compiled Functions

Mathematica can 'compile' certain functions for efficiency. This is not compilation into assembler language, but rather a strong typing of an expression such that intermediate data types do not have to be determined dynamically.

fun[x_] := Exp[-((x - 3)^2/5)]
cfun = Compile[{x}, Exp[-((x - 3)^2/5)]]
time[f_] := Timing[Table[f[1.2], {10^5}]][[1]]
time[fun] 2.4 Second
time[cfun] 0.43 Second

Compile is implicit in many numerical functions, e.g. in Plot.

In a similar manner, Dispatch hashes long lists of rules beforehand, to make the actual substitution faster.

Blocks and Modules

Block implements Dynamical Scoping

A local variable is known everywhere, but only for as long as the block executes ("temporal localization").

Module implements Lexical Scoping

A local variable is known only in the block it is defined in ("spatial localization"). This is how scoping works in most high-level languages.

DownValues and UpValues

Definitions are usually assigned to the symbol being defined: this is called DownValue.

For seldomly used definitions, it is better to assign the definition to the next lower level: this is an UpValue.

This is better than assigning to Plus directly, because Plus is a very common operation.

In other words, Mathematica "looks" one level inside each object when working off transformations.

Mathematica knows some functions to be Output Forms. These are used to format output, but don't "stick" to the result:

{{1, 2}, {3, 4}}//MatrixForm $= \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix}$



Head [%] Car List - not MatrixForm

Some important output forms: InputForm, FullForm, Shallow, MatrixForm, TableForm, TeXForm, CForm, FortranForm.

TeXForm[alpha/(4 Pi)] TeXForm[alpha/(4 CForm[alpha/(4 Pi)] @ alpha/(4.*Pi)

FullForm[alpha/(4 Pi)]

Times[Rational[1, 4], alpha, Power[Pi, -1]]

MathLink

The MathLink API connects Mathematica with external C/C++ programs (and vice versa). J/Link does the same for Java.

```
:Begin:
:Function:
               copysign
               CopySign[x_?NumberQ, s_?NumberQ]
:Pattern:
:Arguments: {N[x], N[s]}
:ArgumentTypes: {Real, Real}
:ReturnType: Real
:End:
#include "mathlink.h"
double copysign(double x, double s) {
 return (s < 0) ? -fabs(x) : fabs(x);
}
int main(int argc, char **argv) {
 return MLMain(argc, argv);
```

In-depth tutorial: http://library.wolfram.com/infocenter/TechNotes/174

T. Hahn, Programming Concepts in Mathematica – p.22

Scripting Mathematica

Efficient batch processing with Mathematica:

Put everything into a script, using sh's Here documents:

```
#! /bin/sh ..... Shell Magic
math << \_EOF_ .... start Here document (note the \)
  << FeynArts'
  << FormCalc'
  top = CreateTopologies[...];
  ...
_EOF_ .... end Here document</pre>
```

Everything between " $<< \tag$ " and "tag" goes to Mathematica as if it were typed from the keyboard.

Note the "\" before tag, it makes the shell pass everything literally to Mathematica, without shell substitutions.

Scripting Mathematica

- Everything contained in one compact shell script, even if it involves several Mathematica sessions.
- Can combine with arbitrary shell programming, e.g. can use command-line arguments efficiently:

```
#! /bin/sh
math -run "arg1=$1" -run "arg2=$2" ... << \END
...
END</pre>
```

• Can easily be run in the background, or combined with utilities such as make.

Debugging hint: -x flag makes shell echo every statement, #! /bin/sh -x

Mathematica Summary

- Mathematica makes it wonderfully easy, even for fairly unskilled users, to manipulate expressions.
- Most functions you will ever need are already built in. Many third-party packages are available at MathSource, http://library.wolfram.com/infocenter/MathSource.
- When using its capabilities (in particular list-oriented programming and pattern matching) right, Mathematica can be very efficient.
 Wrong: FullSimplify[veryLongExpression].
- Mathematica is a general-purpose system, i.e. convenient to use, but not ideal for everything.
 For example, in numerical functions, Mathematica usually selects the algorithm automatically, which may or may not be a good thing.

Mathematica vs. FORM

Mathematica



- Much built-in knowledge,
- Slow if used indiscriminately,
- Very general,
- GUI, add-on packages...

FORM



- Limited mathematical knowledge,
- Fast also on large problems (if known how to handle),
- Optimized for certain classes of problems,
- Batch program (edit-run cycle).

Reference Books, Formula Collections

- V.I. Borodulin et al. CORE (Compendium of Relations) hep-ph/9507456.
- Herbert Pietschmann
 Formulae and Results in Weak Interactions Springer (Austria) 2nd ed., 1983.
- Andrei Grozin
 Using REDUCE in High-Energy Physics
 Cambridge University Press, 1997.

Antisymmetric Tensor

The Antisymmetric Tensor in *n* dimensions is denoted by $\varepsilon_{i_1i_2...i_n}$. You can think of it as a matrix-like object which has either -1, 0, or 1 at each position.

For example, the Determinant of a matrix, being a completely antisymmetric object, can be written with the ε -tensor:

$$\det A = \sum_{i_1,\dots,i_n=1}^n \varepsilon_{i_1i_2\dots i_n} A_{i_11} A_{i_22} \cdots A_{i_nn}$$

In practice, the ε -tensor is usually contracted, e.g. with vectors. We will adopt the following notation to avoid dummy indices:

$$\varepsilon_{\mu\nu\rho\sigma}p^{\mu}q^{\nu}r^{\rho}s^{\sigma}=\varepsilon(p,q,r,s).$$

Epsilon tensor in Mathematica

(* for actual vectors, this evaluates to the determinant, but take care of the signs from the g_{mu nu}s *)

Eps[args__List] := I (-1)^Length[{args}] Det[{args}]

(* implement linearity: *)

Eps[a___, p_Plus, b___] := Eps[a, #, b]&/@ p

 $Eps[a_{, n}, n_{v}] = n Eps[a, r, b]$

(* otherwise sort the arguments into canonical order: *)
Eps[args__] := Signature[{args}] Eps@@ Sort[{args}] /;
!OrderedQ[{args}]

Abbreviationing

One of the most powerful tricks to both reduce the size of an expression and reveal its structure is to substitute subexpressions by new variables.

The essential function here is Unique with which new symbols are introduced. For example,

Unique["test"]

generates e.g. the symbol test1, which is guaranteed not to be in use so far.

The Module function which implements lexical scoping in fact uses Unique to rename the symbols internally because Mathematica can really do dynamical scoping only.

Abbreviations in Mathematica

(* the main abbreviationing function *)

```
$AbbrPrefix = "c";
abbr[expr_] := abbr[expr] = Unique[$AbbrPrefix]
```

```
(* apply abbr e.g. like this: *)
```

```
Structure[expr_, x_] := Collect[expr, x, abbr]
```

(* get list of abbreviations introduced so far *)

AbbrList[] := Cases[DownValues[abbr], _[_[_[f_]], s_Symbol] -> (s -> f)]

Restore[expr_] := expr /. AbbrList[]

T. Hahn, Programming Concepts in Mathematica – p.31

Colour Structures

In Feynman diagrams four type of Colour structures appear:

a llll Representation $\sim T^a_{ii} = ext{SUNT} [a, i, j]$ Natural $T \sim T^a_{ij} T^a_{k\ell} = \texttt{SUNTSum}[i, j, k, \ell]$

a eccepter b elle c $a \sim f^{abc} = ext{SUNF} [a, b, c]$ Keo dioint soor ellele d b ~66666 $\sim f^{abx} f^{xcd} = ext{SUNF}[a, b, c, d]$

Unified Notation

The SUNF's can be converted to SUNT's via

$$f^{abc} = 2\mathrm{i} \left[\mathrm{Tr}(T^c T^b T^a) - \mathrm{Tr}(T^a T^b T^c) \right]$$

We can now represent all colour objects by just SUNT:

• SUNT [
$$i$$
 , j] $= \delta_{ij}$

- SUNT [a, b, ..., \overline{i} , j] = $(T^a T^b \cdots)_{ij}$
- SUNT [a,b,...,0,0] = $\operatorname{Tr}(T^a T^b \cdots)$

This notation again avoids unnecessary dummy indices. (Mainly namespace problem.)

For purposes such as the "large- N_c limit" people like to use SU(N) rather than an explicit SU(3).

T. Hahn, Programming Concepts in Mathematica - p.33

Fierz Identities

The Fierz Identities relate expressions with different orderings of external particles. The Fierz identities essentially express completeness of the underlying matrix space.

They were originally found by Markus Fierz in the context of Dirac spinors, but can be generalized to any finite-dimensional matrix space [hep-ph/0412245].

For SU(N) (colour) reordering, we need

$$T^a_{ij}T^a_{k\ell} = \frac{1}{2}\left(\delta_{i\ell}\delta_{kj} - \frac{1}{N}\delta_{ij}\delta_{k\ell}\right).$$

T. Hahn, Programming Concepts in Mathematica – p.34

Cvitanovich Algorithm

For an Amplitude:

- convert all colour structures to (generalized) SUNT objects,
- simplify as much as possible, i.e. use the Fierz identity on all internal gluon lines.

For a Squared Amplitude:

• use the Fierz identity for SU(N) to get rid of all SUNT objects.

For "hand" calculations, a pictorial version of this algorithm exists in the literature.

Translation to Colour-Chain Notation

In colour-chain notation we can distinguish two cases:

a) Contraction of different chains:

$$\langle A | T^a | B \rangle \langle C | T^a | D \rangle = \frac{1}{2} \left(\langle A | D \rangle \langle C | B \rangle - \frac{1}{N} \langle A | B \rangle \langle C | D \rangle \right),$$

b) Contraction on the same chain:

$$\langle A | T^a | B | T^a | C \rangle = rac{1}{2} \left(\langle A | C \rangle \operatorname{Tr} B - rac{1}{N} \langle A | B | C \rangle \right).$$

T. Hahn, Programming Concepts in Mathematica – p.36

Colour Algebra in Mathematica

(* in-chain version of the Fierz identity *)

sunT[t1___, a_Symbol, t2___, a_, t3___, i_, j_] :=
 (sunT[t1, t3, i, j] (sunT[t2, #, #]&[Unique["c"]]) sunT[t1, t2, t3, i, j]/SUNN)/2

(* across-chain version of the Fierz identity *)

```
sunT[t1___, a_Symbol, t2___, i_, j_] *
sunT[t3___, a_, t4___, k_, l_] ^:=
  (sunT[t1, t4, i, 1] sunT[t3, t2, k, j] -
      sunT[t1, t2, i, j] sunT[t3, t4, k, 1]/SUNN)/2
```

(* apply e.g. like this: *)

ColourSimplify[expr_] := expr /. SUNT -> sunT

Fermion Trace

Leaving apart problems due to γ_5 in d dimensions, we have as the main algorithm for the 4d case:

$$\operatorname{Tr} \gamma_{\mu} \gamma_{\nu} \gamma_{\rho} \gamma_{\sigma} \cdots = + g_{\mu\nu} \operatorname{Tr} \gamma_{\rho} \gamma_{\sigma} \cdots \\ - g_{\mu\rho} \operatorname{Tr} \gamma_{\nu} \gamma_{\sigma} \cdots \\ + g_{\mu\sigma} \operatorname{Tr} \gamma_{\nu} \gamma_{\rho} \cdots$$

This algorithm is recursive in nature, and we are ultimately left with

$$\operatorname{Tr} 1 = 4.$$

(Note that this 4 is not the space-time dimension, but the dimension of spinor space.)

T. Hahn, Programming Concepts in Mathematica – p.38

Fermion Trace in Mathematica

(* pick out one index, mu, at a time *)

```
Trace4[mu_, g__] :=
Block[ {Trace4, s = -1},
    Plus@@ MapIndexed[
        ((s = -s) Pair[mu, #1] Drop[Trace4[g], #2])&,
        {g} ] ]
```

(* the unit trace *)

Trace4[] = 4

T. Hahn, Programming Concepts in Mathematica - p.39