







# Optimizations for HEP Jobs on HPC with XRootD Caching

Robin Hofsaess for the GridKa R&D team





# **Future German HEP Computing Strategy**



#### Plan:

- Transition from dedicated T2 resources at universities to shares on national HPC centers within the <u>NHR computing</u> compound
- Storage at the Helmholtz centers KIT and DESY
- Support from the current T2 groups

## **Challenges:**

- No influence on the hardware setup
- From admin to a user among many
- Different compute models
- WLCG pledges

#### Why HPC?

- Use resources that are available anyway
- Increased efficiency and sustainability
- Enhanced collaboration

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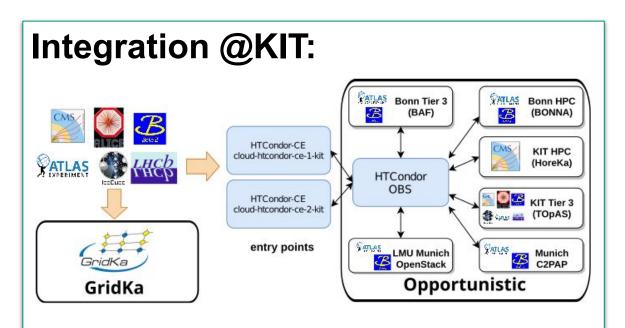
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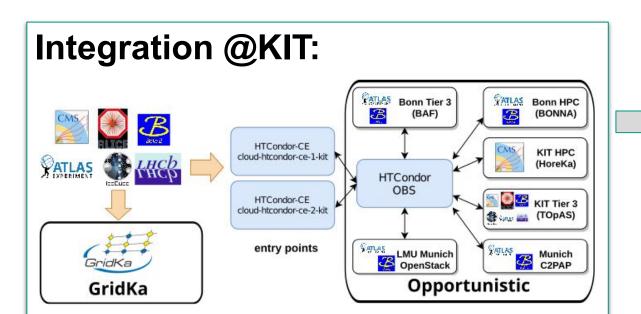
A reliable and efficient integration of HPC requires R&D and testing!





HoreKa is integrated **opportunistically** with **COBaID/TARDIS** since more than **three** years





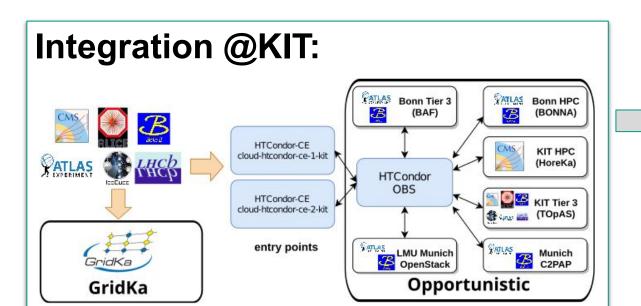
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#### **Observations**

In comparison to T1/T3:

- Increased failure rate
- lower CPU efficiency



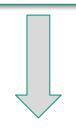


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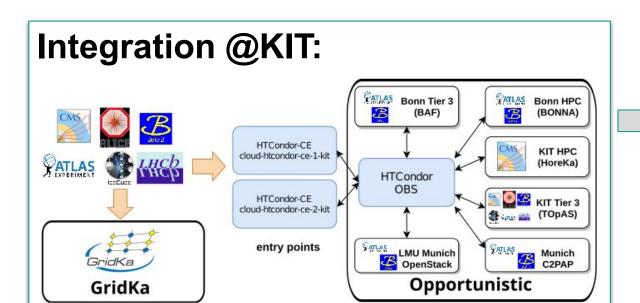
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Why and how to find out?

ACAT Contribution





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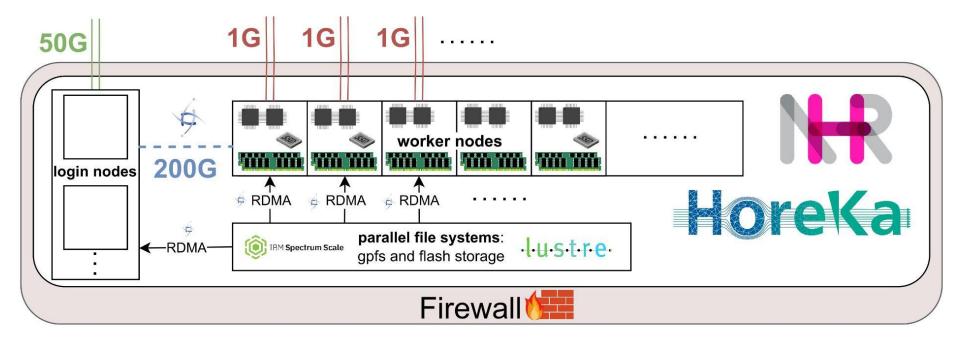
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The reasons can be manifold and require a close look at each individual HPC center!

## HoreKa Overview

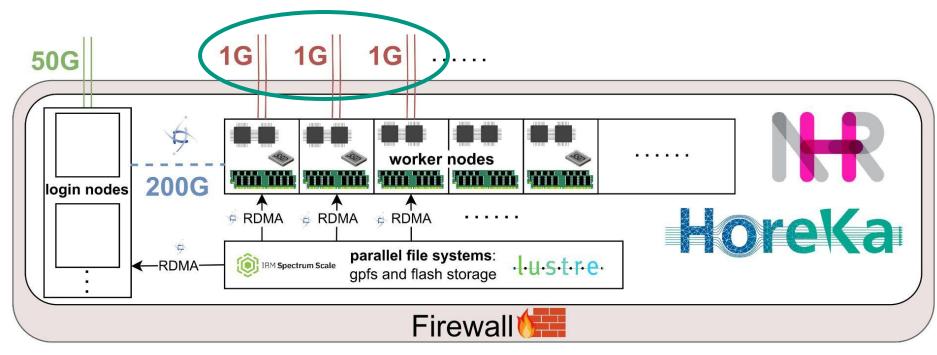




- Separate test projects per VO
- A CMS test project got approved recently
- Additionally: backfilling of WLCG jobs dependent of cluster utilization

## HoreKa Overview



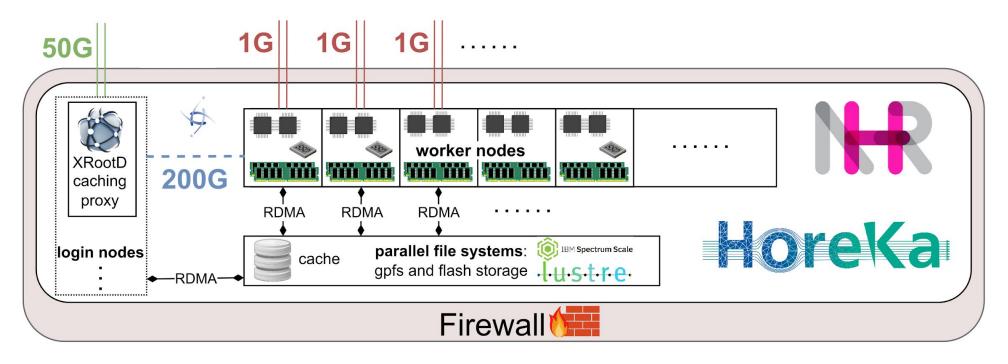


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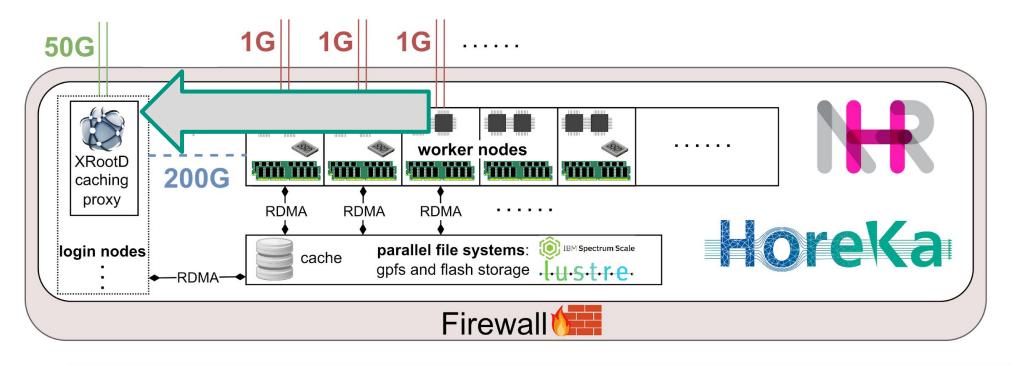
Bandwidth limitations were identified as the main bottleneck for data intensive workflows





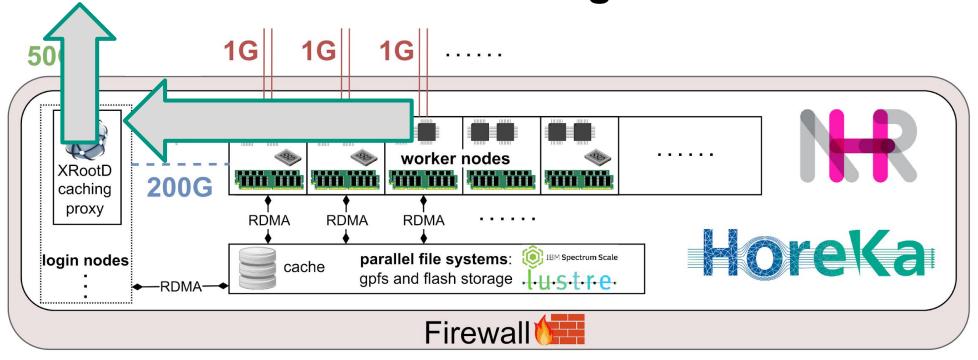






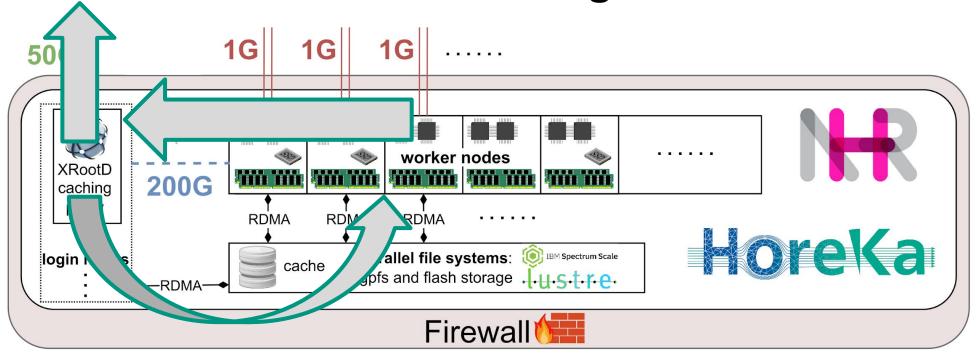
- The traffic is proxied over a login node via XRootD (50G connectivity)
- Remote data can be cached on the fly on the parallel filesystem
- Conceptually, this corresponds to a buffer





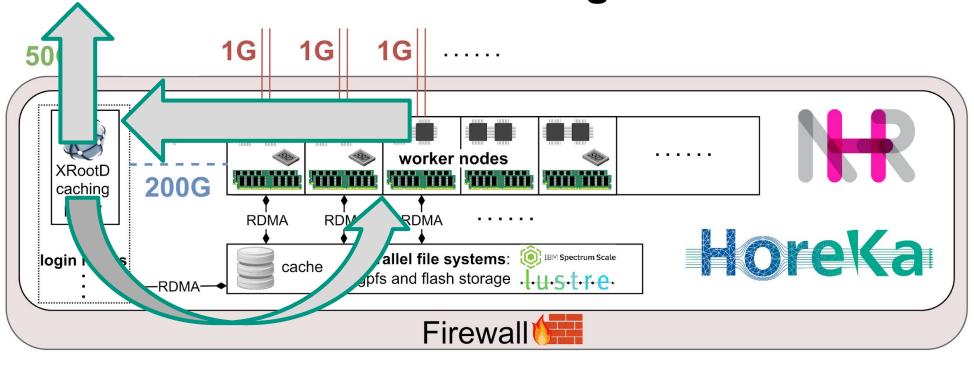
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Now deployed in production for further testing!

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## **Advantages of our Concept**

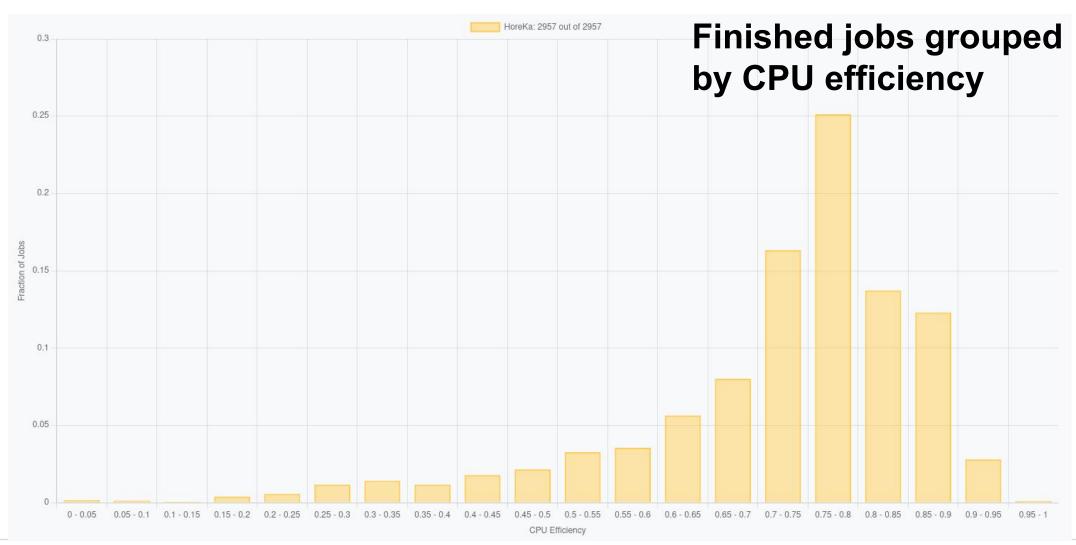


## **Faster Connectivity:**

The **prefetching** of data over the login node accelerates the remote data access.

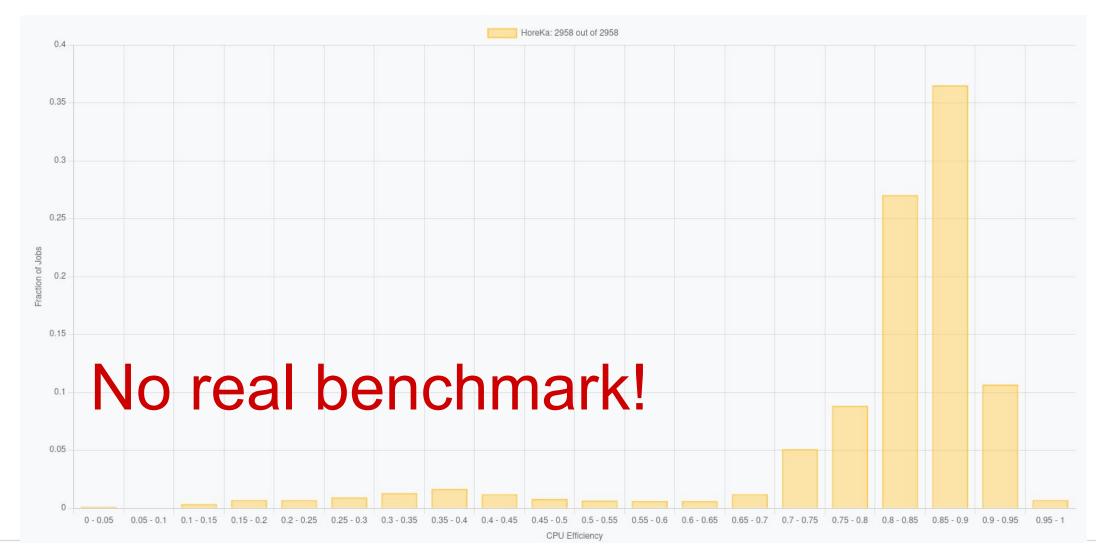
# **Preliminary Results: No Caching Proxy**





# **Preliminary Results: With Caching Proxy**





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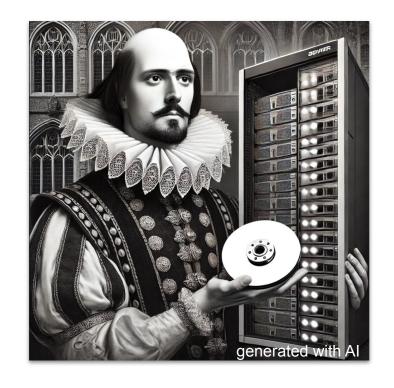
## Caching:

Cache hits may increase the transfer speed a lot (up to several GB/s!)

## To cache or not to cache? That is the question!



- This is a tough decision:
  - Can be very useful with a decent cache hit rate:
    - Depends on job mix, data sets, cache size, etc
    - E.g.: very positive experiences with caching of user analysis data in Spain (<u>ACAT poster</u>)
  - From our experience on HPC:
     bleeding edge and still error-prone
  - BUT: Very fruitful collaboration with the XRootD team for improvements [1, 2]



- Currently hard to decide
- Further testing and optimizations are necessary and planned
- We are working on additional studies and benchmark tests

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#### **Additional Benefits:**

- Adds monitoring (very valuable for site operation!)
- Opens <u>new ways</u> of integrating HPC centers

## **Conclusion and Outlook**



- We have developed and deployed a PoC at HoreKa
- The setup runs very smooth and is easily adaptable
- Failure rate and CPU efficiency overall improved and are comparable with limited jobmix
- Also beneficial without a data access bottleneck
- The XRootD proxy cache helps a lot in terms of site operation by providing logs/monitoring data

The prefetching over the login node effectively mitigates remote data access bottlenecks

Caching requires more studies but may be useful in a future German data lake infrastructure with further adaptions

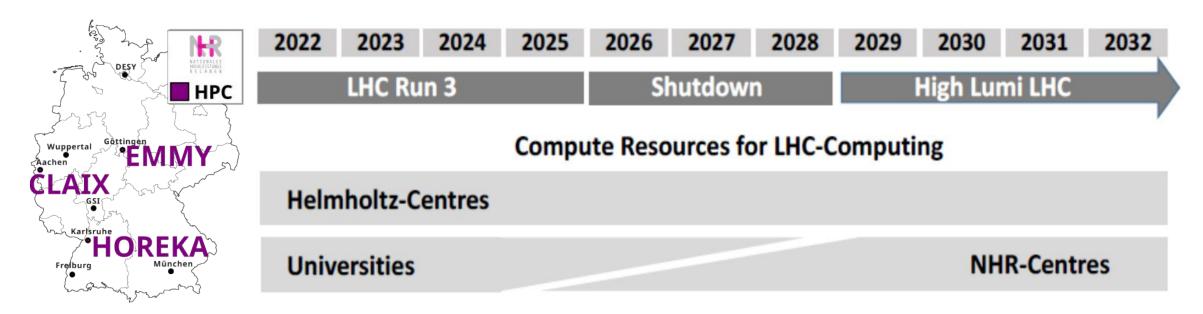
Our concept adds further benefits and possibilities to optimize the HPC integration besides the transfer speed increasement for data access



# THANKS!

# **German HEP Computing Strategy**



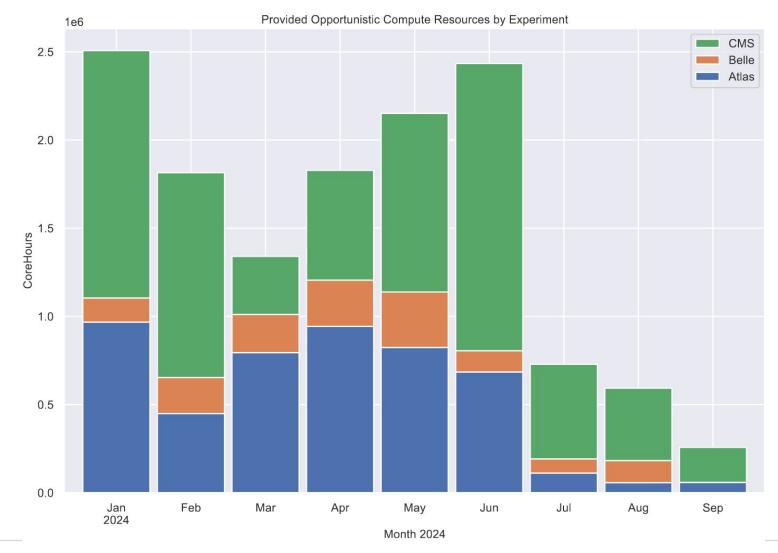


- Transition from dedicated T2 resources at universities to shares on national HPC centers within the <a href="NHR computing">NHR computing</a> compound
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strategy paper

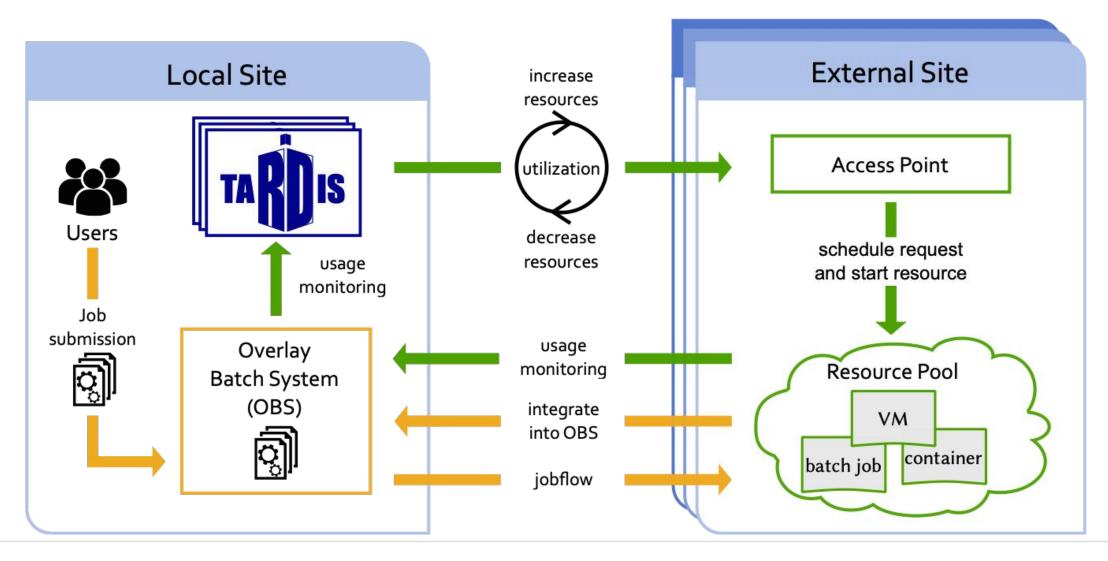






## **COBaID/TARDIS**

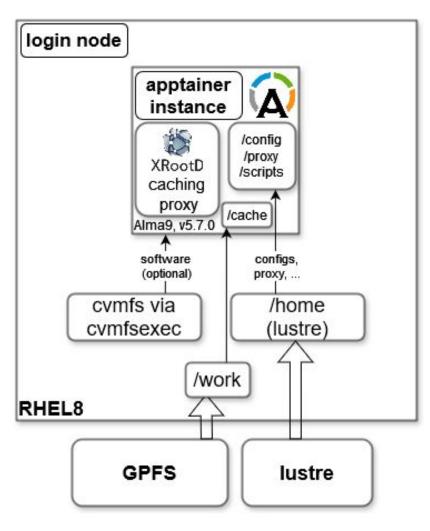




## Setup and Configuration: Overview



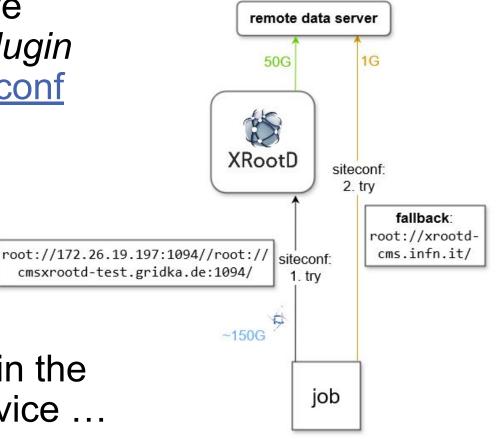
- Host: RHEL8
- usernamespaces, CGroups v2, systemd user services
- Currently running: XRootD v5.7.1 as Alma9 apptainer instance (image bootstrapped from docker)
- In principle up to 76c and 500GB RAM, 50G WAN
  - But shared with other users (limitation via apptainer instance with CGv2 possible)
  - Usage: 32t, 64GB memory
- 250T quota on gpfs (via IB)



# **Setup and Configuration: Connection**



- For enabling the proxy for transfers, we currently do not use the XrdCIProxyPlugin
- Instead, the proxy is added to the <u>siteconf</u> directly
- Advantages:
  - We can use the intended fallback mechanism if smth fails
  - It is only enabled for file reads
- Disadvantages:
  - Changes require always a change in the repo and a full shutdown of the service ...



# **Setup and Configuration: RDMA**



- We currently don't use RDMA natively, but <u>IPoIB</u> for the transfers
- Currently, only IPv4 is possible (no link local v6 addresses in xrootd)
- The cache is also mounted via IPoIB (RDMA)
- We got some complaints from the GPFS team because of the many, many IOPS
- Reason: small blocksizes when caching is enabled
- ideally (FS PoV): pfc.blocksize == FS blocksize, or in general: as big as possible
- dca:
  - Tested, but problematic with containerization
  - Dependent of the campaigns/datasets, e.g. premix rarely completely cached
     in production rather pointless
  - Would be very useful, if made possible for partially cached files if possible.

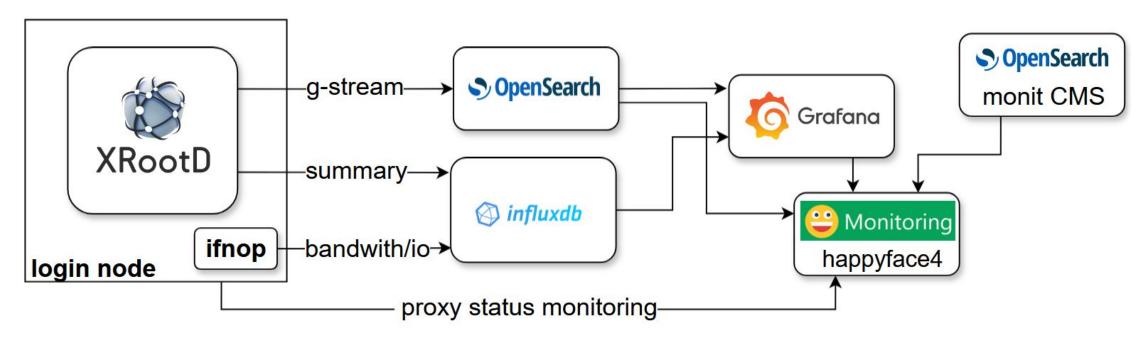
## **Job Mix**



- Currently, we only run a subset of the full job mix
- This works well, even for more data intensive WFs
- A full replacement will require the full mix, including Analysis
- Caching may be useful for a subset of the job mix (e.g. not RAW)
  - TBD with more monit/benchmarks
- Our proposal:
  - Even the full replacement should not run jobs like *Merge*, as they are just too inefficient on the expensive HPC hardware
  - To achieve the best possible efficiency, it is crucial to consider that

# **Monitoring**



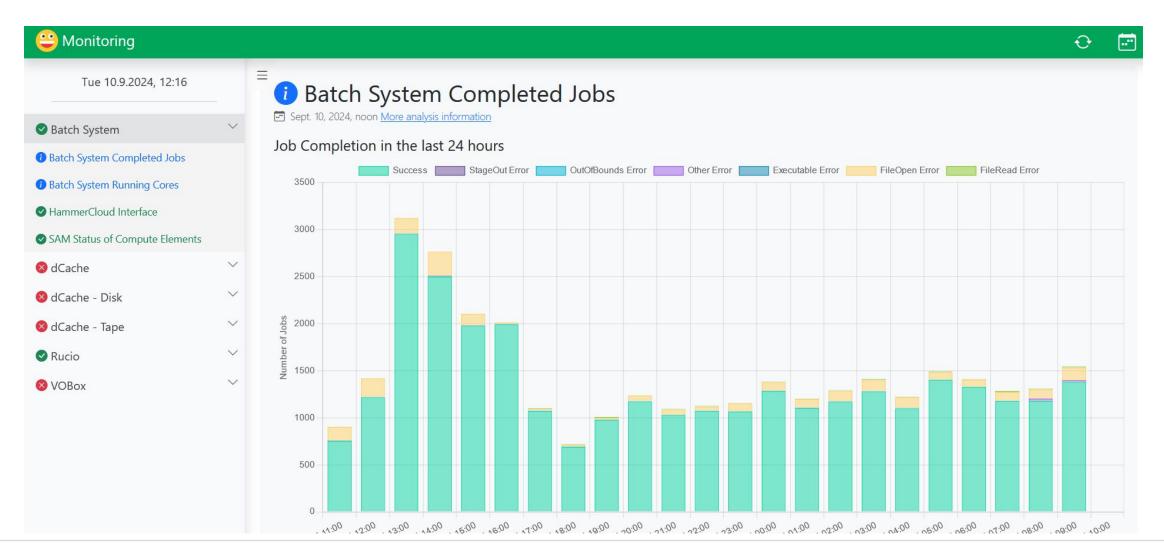


- Many additional monitoring capabilities thanks to XRootD
- Especially cache summary monitoring would be helpful (on the way)
- We collect everything and unify it in our meta-monitoring: <a href="HappyFace4"><u>HappyFace4</u></a>
- I learned a lot about CMS and created some tools: <u>I</u>, <u>II</u>

Note: not yet fully in production/public







## **Benchmark Mechanism**



- We are working on requestable benchmarking and debugging WFs
  - Will help for debugging, as it provides a testing scenario that is well understood and can be more verbose (in comparison to production jobs – e.g. no debug flags etc)
  - Will provide a more comparable benchmark mechanism for site comparison e.g. to answer:
    - "Did my recent changes actually work or am I just lucky with the scheduling"
- Oriented on the ARM RelVal workflows

## Improvements and Plans



- Further improvements:
  - Caching improvements: E.g. only cache certain sites

•

- Plans:
  - Make the most out of all monitoring
  - Switch to tokens
  - Investigate feasibility and caching efficiency
  - Develop a benchmark mechanism
  - Further data analysis for comparing HPC with "normal" grid sites
  - Documentation and publishing coming soon