Status report

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Current Work

- Integration of Lustre HSM and Gekkofs
- Providing caching by using ad-hoc file systems
- Compatibility of HPC applications and ad-hoc file systems
- Analysis of BESIII & LOFAR workloads on MOGON II and MOGON-NHR



HPC APPLICATION ANALYSIS

- HPC I/O Analysis
- Released at : <u>https://hpcioanalysis.zdv.uni-mainz.de/</u>
- Capturing I/O behavior with DARSHAN
- Analyzing I/O behavior
- Ongoing applications:

BESIII

LOFAR

AD-HOC FILE SYSTEMS

- Employing unused node-local storages like SSDs and NVMs
- Distributing data and metadata across node-local storage
- Parallel read/write on multiple nodes
- Hiding slow performance of backend storage
- Reducing I/O bottlenecks

AD-HOC FILESYSTEMS LIMITATIONS

- User responsibilities with ad-hoc file systems
- Transferring data to node-local storage
- Separated namespace with back-end storage
- Risk of data inconsistency
- Node-local storage integration into the storage hierarchy is still a problem
- Requiring restart of the job when compute node is restarting

EXISTING SOLUTION - LPCC

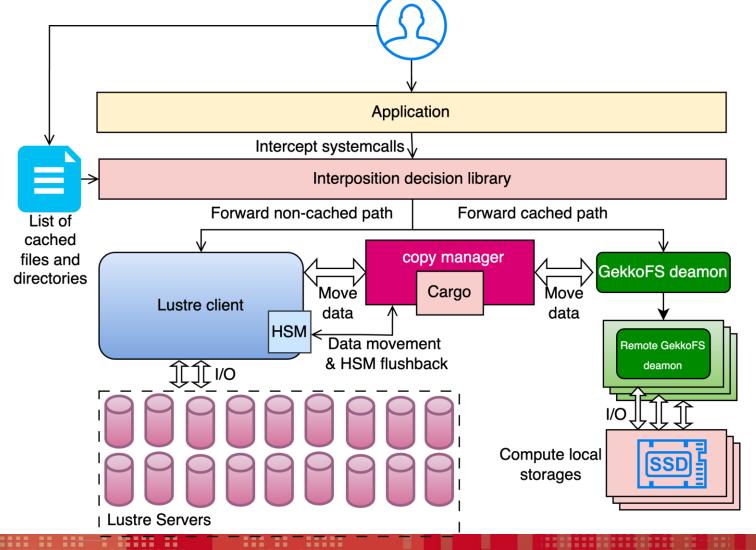
- Offers client cache to user
- Employs HSM from Lustre to guarantee consistency
- RW mode
 - Transfers files which accessed exclusively by one node
- Read mode
 - Transfers a version of read-only data to each node
- No single namespace



LUSTRE HSM INTEGRATION WITH GEKKOFS

- Providing a single namespace
- Guarantee data consistency
- Providing a cache for HPC application with node local storage
- Integration of node local storage into the storage hierarchy
- Minimize user interaction with transparent design
- Automatic data transfer

INTEGRATION LUSTRE HSM AND GEKKOFS



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DESIGN FEATURES

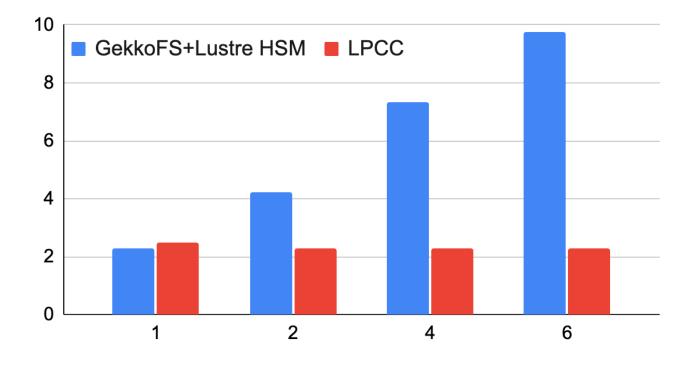
- User can define files and directories path for caching
- Using interception library to dispatch requests between Lustre and Gekkofs
- Using Cargo as a copy tool
- Using Lustre HSM to offer single namespace
- HSM is storing the status of a file (existing on Lustre, moved to Gekkofs)
- HSM triggers a file flush back from Gekkofs to Lustre in the case of conflicting access
- **Read-only mode**



READ PERFORMANCE – HSM INTEGRATED

3andwidth (GB/s)

- IOR test
- Read each file 10 times
- One file per IOR
- IOR transfersize = 1 MB
- 32 processes per each node
- 4 GB per process
- Overhead: an extra read from Lustre and an extra write on NVM
- LPCC can not cache shared data between nodes

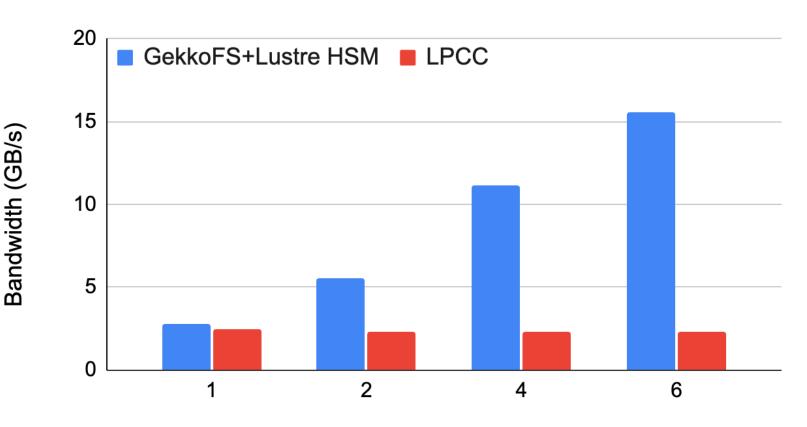


Number of nodes

Lustre min bandwidth = 1071MiB/s Lustre max bandwidth =3059 MiB/s Lustre average bandwidth = 2620 MiB/s

WRITE PERFORMANCE- HSM INTEGRATED

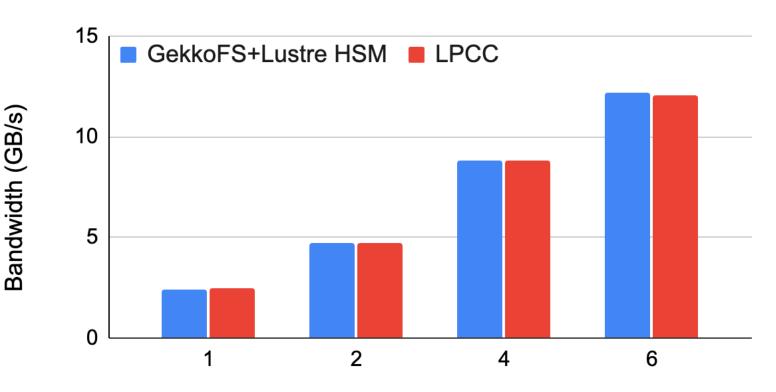
- IOR test
- Read file 10 times
- One file per IOR
- IOR transfersize = 1 MB
- 32 processes per each node
- 4 GB per process
- Creates file directly on NVM
- No additional overhead



Number of nodes

READ PERFORMANCE – HSM INTEGRATED

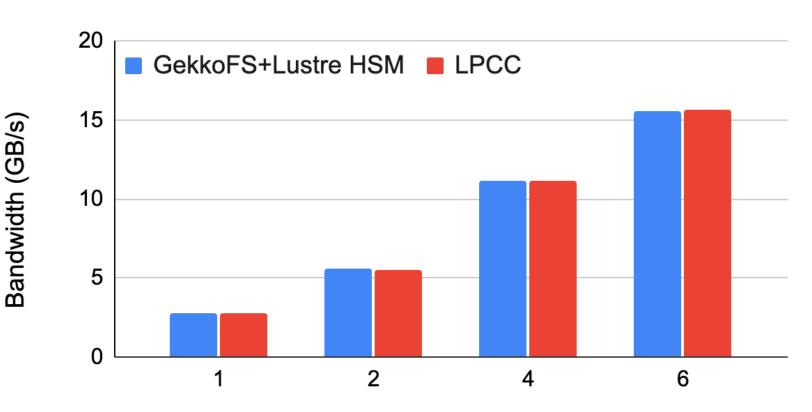
- IOR test
- Read file 10 times
- One file per process
- IOR transfersize = 1 MB
- 32 processes per each node
- 4 GB file per node
- Overhead: an extra read from Lustre and an extra write on NVI
- LPCC can cache unique files per node



Number of nodes

WRITE PERFORMANCE- HSM INTEGRATED

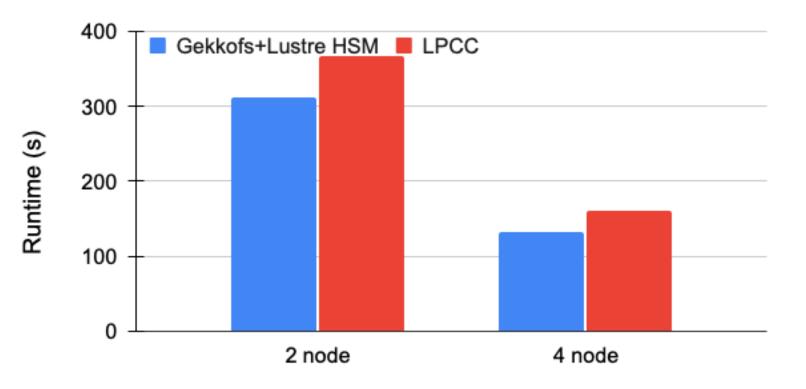
- IOR test
- Read file 10 times
- One file per process
- IOR transfersize = 1 MB
- 32 processes per each node
- 4 GB file per node
- Creates file directly on NVM
- No additional overhead



Number of nodes

NEK5000

- With 50 steps
- 16 processes per node
- Writes 10 output file
- Each file less than 700 MB
- Stage-in time less than 1 second

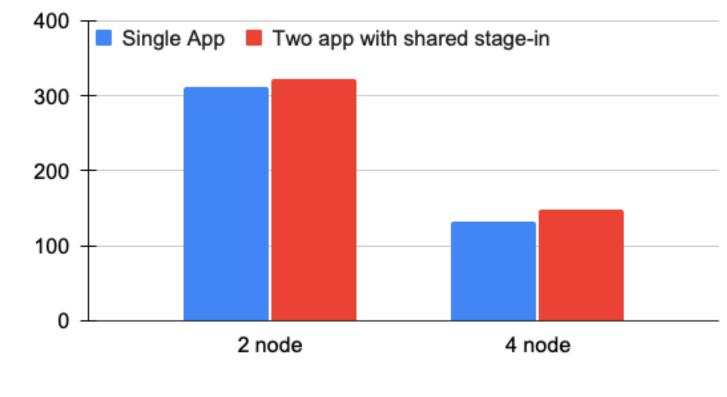


Number of nodes



NEK5000

- Running two Nek5000 applications
- Shared stage-in dataset
- Each application stage-in a copy of data to node-local storage
- Writes 10 output file
- Reporting the worst runtime between two applications
- Each application separately working on their dataset



Number of nodes



Runtime (s)



Done:

- Adapting Darshan to trace BESIII
- Analyzing I/O behavior in small-scale configurations
 Ongoing/todo:
- Adapting GekkoFS to run inside BESIII container
- Running BESIII with larger scale configurations
- Providing I/O analysis and possible improvements
- Optimizing performance of BESIII by using GekkoFS



Done:

- Analysed I/O performance in wsclean stage
- Evaluating performance improvement of using local scratch device over parallel file systems

Ongoing:

- Adapting GekkoFS to run LOFAR
- Reducing LOFAR runtime by asynchronously staging data





- Intelligent data staging
- File system interference with applications and NVMe over Fabrics (NVMe-oF)





- Adding node-local SSDs to storage hierarchy
- Transparent solution for using ad-hoc storage
- Analyzing BESIII and LOFAR on MOGON II & MOGON-NHR
- Compatibility of HPC Application
- Analyzing HPC traces and applications

THANK YOU

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Gitlab-Repo: https://storage.bsc.es/gitlab/hpc/gekkofs/



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