

On Solutions of Recurrences I

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Outline

- 1 Summation and recurrences
- 2 What is a solution?
- 3 C-recursive sequences
- 4 P-recursive sequences
- 5 Recurrence operators
- 6 Difference rings
- 7 Solutions of special types

Summation and recurrences

$$s_n = \sum_{k=0}^{n-1} t_k$$



$$s_{n+1} - s_n = t_n; \quad s_0 = 0$$



$$t_n s_{n+2} - (t_n + t_{n+1}) s_{n+1} + t_{n+1} s_n = 0; \quad s_0 = 0, s_1 = t_0$$

Representation of sequences

- *recursive*: a_0, a_1, \dots, a_{d-1} given,

$$a_n = F(a_{n-1}, a_{n-2}, \dots, a_0, n) \quad \text{for } n \geq d$$

- *explicit*: $a_n = f(n)$ for $n \geq 0$

- *by generating function*:

$$G_a(x) = \sum_{n=0}^{\infty} a_n x^n$$

What is a solution?

Example

- *recursive*: $a_0 = 0, a_1 = 1,$

$$a_n = 2a_{n-1} - a_{n-2} \quad \text{for } n \geq 2$$

- *explicit*: $a_n = n \quad \text{for } n \geq 0$

- *by generating function*:

$$G_a(x) = \frac{x}{(1-x)^2}$$

Notation:

K ... algebraically closed field of characteristic 0

Definition

A sequence $\langle a_n \rangle_{n=0}^{\infty} \in K^{\mathbb{N}}$ is *C-recursive* if there are $d \in \mathbb{N}$ and constants $c_1, c_2, \dots, c_d \in K$, $c_d \neq 0$, such that

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_d a_{n-d}$$

for all $n \geq d$.

C-recursive sequences

Let $\langle a_n \rangle_{n=0}^{\infty} \in K^{\mathbb{N}}$ and $G_a(x) = \sum_{n=0}^{\infty} a_n x^n$.

Theorem

The following are equivalent:

- 1 $\langle a_n \rangle_{n=0}^{\infty}$ is C-recursive,
- 2 $a_n = \sum_{i=1}^r P_i(n) \alpha_i^n$ where $P_i \in K[n]$ and $\alpha_i \in K$,
- 3 $G_a(x) = \frac{P(x)}{Q(x)}$ where $P, Q \in K[x]$, $\deg P < \deg Q$ and $Q(0) \neq 0$.

Theorem

*C-recursive sequences are closed under the following **binary operations** $(a, b) \mapsto c$:*

- 1** *addition:* $c_n = a_n + b_n$
- 2** *Hadamard (termwise) multiplication:* $c_n = a_n b_n$
- 3** *Cauchy multiplication (convolution):* $c_n = \sum_{i=0}^n a_i b_{n-i}$
- 4** *interlacing:* $\langle c_0, c_1, c_2, c_3, \dots \rangle = \langle a_0, b_0, a_1, b_1, \dots \rangle$

Remark

These operations extend naturally to an arbitrary nonzero finite number of operands.

Theorem

*C-recursive sequences are closed under the following **unary operations** $a \mapsto c$:*

- 1 *scalar multiplication:* $c_n = \lambda a_n$ ($\lambda \in K$)
- 2 *(left) shift:* $c_n = a_{n+1}$
- 3 *indefinite summation:* $c_n = \sum_{k=0}^n a_k$
- 4 *multisection:* $c_n = a_{kn+r}$ ($k \in \mathbb{N}$, $0 \leq r < k$)

Example

- $a_n = n + 1$ is C-recursive,
- $a_n^{-1} = \frac{1}{n + 1}$ is **not** C-recursive.

Question: When are a and $1/a$ are both C-recursive?

C-recursive sequences

Definition

A sequence $a \in K^{\mathbb{N}}$ is *geometric* if $a_0 \neq 0$ and $\exists q \in K^*$:

$$a_n = q a_{n-1}$$

for all $n \geq 1$ (equivalently: $a_n = a_0 q^n$ for all $n \geq 0$).

Observation

a geometric $\implies a^{-1}$ geometric

Theorem

(Larson & Taft, 1990) Sequences a and $1/a$ both C-recursive
 $\iff a$ is the interlacing of one or more geometric sequences.

Definition

A sequence $\langle a_n \rangle_{n=0}^{\infty} \in K^{\mathbb{N}}$ is *P-recursive* if there are $d \in \mathbb{N}$ and polynomials $p_0, p_1, \dots, p_d \in K[n]$, $p_d \neq 0$, such that

$$p_d(n)a_{n+d} + p_{d-1}(n)a_{n+d-1} + \dots + p_0(n)a_n = 0$$

for all $n \geq 0$.

Definition

A f.p.s. $f(x) \in K[[x]]$ is *D-finite* if there are $d \in \mathbb{N}$ and polynomials $q_0, q_1, \dots, q_d \in K[x]$, $q_d \neq 0$, such that

$$q_d(x)f^{(d)}(x) + q_{d-1}(x)f^{(d-1)}(x) + \dots + q_0(x)f(x) = 0.$$

Let $\langle a_n \rangle_{n=0}^{\infty} \in K^{\mathbb{N}}$ and $G_a(x) = \sum_{n=0}^{\infty} a_n x^n$.

Theorem

The following are equivalent:

- 1 $\langle a_n \rangle_{n=0}^{\infty}$ is P-recursive,
- 2 $G_a(x)$ is D-finite.

Theorem

P-recursive sequences are closed under the following operations:

- 1 *addition*
- 2 *Hadamard multiplication*
- 3 *Cauchy multiplication*
- 4 *interlacing*
- 5 *scalar multiplication*
- 6 *shift*
- 7 *indefinite summation*
- 8 *multisection*

Example

- $a_n = 2^n + 1$ is P-recursive (even C-recursive),
- $b_n := a_n^{-1} = \frac{1}{2^n + 1}$ is **not** P-recursive.

Sketch of proof:

$$G_b(x) = \sum_{n=0}^{\infty} b_n x^n = \sum_{n=0}^{\infty} \frac{x^n}{2^n + 1}$$

radius of convergence:

$$r = \frac{1}{\limsup_{n \rightarrow \infty} \sqrt[n]{|b_n|}} = \lim_{n \rightarrow \infty} \sqrt[n]{2^n + 1} = 2$$

$$G_b(2x) = \sum_{n=0}^{\infty} \frac{2^n}{2^n + 1} x^n = \sum_{n=0}^{\infty} \left(1 - \frac{1}{2^n + 1}\right) x^n$$
$$G_b(2x) = \frac{1}{1-x} - G_b(x) \quad (1)$$

$x = 1$: $\frac{1}{1-x}$ singular, G_b regular $\implies G_b$ singular at $x = 2$

$x = 2$: $\frac{1}{1-x}$ regular, G_b singular $\implies G_b$ singular at $x = 4$

$x = 4$: $\frac{1}{1-x}$ regular, G_b singular $\implies G_b$ singular at $x = 8$

...

P-recursive sequences

By induction on k : $G_b(x)$ singular at $x = 2^k$ for all $k \in \mathbb{N}$

$\implies G_b$ not D-finite

$\implies b$ not P-recursive □

Question: When are a and $1/a$ are both P-recursive?

Definition

A sequence $a \in K^{\mathbb{N}}$ is *hypergeometric* if:

1 $\exists N \in \mathbb{N} : a_n \neq 0$ for all $n \geq N$,

2 $\exists p, q \in K[n] \setminus \{0\}$:

$$p(n) a_{n+1} + q(n) a_n = 0$$

for all $n \geq 0$.

Observation

a hypergeometric $\implies a^{-1}$ hypergeometric

Theorem

(Singer, 1997) Sequences a and $1/a$ both P-recursive $\iff a$ is the interlacing of one or more hypergeometric sequences.

Notation:

$\mathcal{H}(K) \dots$ hypergeometric sequences in $K^{\mathbb{N}}$

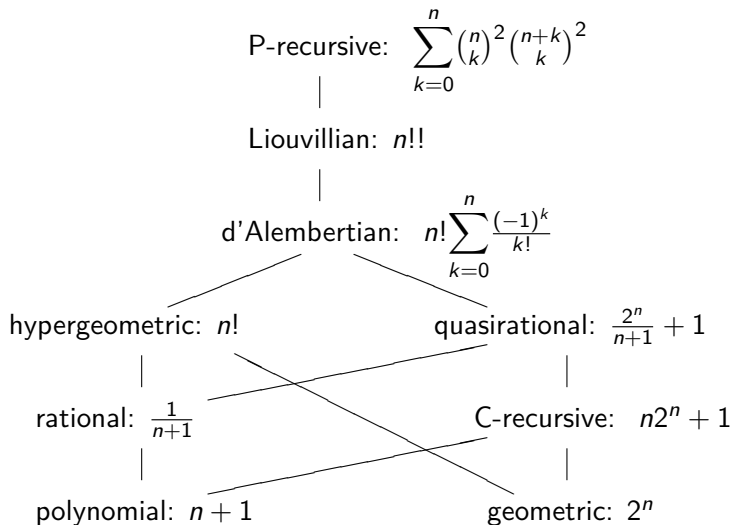
Solving linear recurrences with polynomial coefficients

Given: $d \in \mathbb{N}$ and $p_0, p_1, \dots, p_d \in K[n]$, $p_d \neq 0$

Find: all *nice* solutions $a \in K^{\mathbb{N}}$ of

$$p_d(n)a_{n+d} + p_{d-1}(n)a_{n+d-1} + \cdots + p_0(n)a_n = 0$$

P-recursive sequences



Recurrence operators

Notation:

$$E : K^{\mathbb{N}} \rightarrow K^{\mathbb{N}} \quad (\text{left}) \textit{ shift operator},$$
$$(Ea)_n = a_{n+1} \quad \text{application of } E$$

$$(E^k a)_n = a_{n+k} \quad (k \in \mathbb{N})$$

Given $d \in \mathbb{N}$ and $p_0, p_1, \dots, p_d \in K[n]$, $p_d \neq 0$:

$$L : K^{\mathbb{N}} \rightarrow K^{\mathbb{N}} \quad \textit{linear recurrence operator},$$

$$L = \sum_{k=0}^d p_k E^k$$

$$(La)_n = \sum_{k=0}^d p_k(n) a_{n+k} \quad \text{application of } L$$

Notation:

$K[n]\langle E \rangle$... algebra of linear recurrence operators
with polynomial coefficients

Commutation rule for *composition* of operators:

$$E \cdot p(n) = p(n+1)E$$

$$\sum_{k=0}^d p_k(n)E^k \cdot \sum_{j=0}^e q_j(n)E^j = \sum_{k=0}^d \sum_{j=0}^e p_k(n)q_j(n+k)E^{j+k}$$

Definition

A *difference ring* is a pair (K, σ) where:

- K is a commutative ring with multiplicative identity,
- $\sigma : K \rightarrow K$ is a ring automorphism.

If, in addition, K is a field then (K, σ) is a *difference field*.

Example

- $(K[x], \sigma)$ with $\sigma x = x + 1$, $\sigma|_K = \text{id}_K$ is a difference ring.
- $(K(x), \sigma)$ with $\sigma x = x + 1$, $\sigma|_K = \text{id}_K$ is a difference field.
- $(K^{\mathbb{N}}, E)$ where $E : \langle a_0, a_1, a_2, \dots \rangle \mapsto \langle a_1, a_2, a_3, \dots \rangle$ is *not* a difference ring.

Difference rings

For $a, b \in K^{\mathbb{N}}$ define:

$$a \sim b \iff \exists N \in \mathbb{N} : \forall n \geq N : a_n = b_n$$

Notation:

$\mathcal{S}(K) = K^{\mathbb{N}} / \sim$... ring of **germs of sequences**

$\sigma : \mathcal{S}(K) \rightarrow \mathcal{S}(K)$... unique automorphism of $\mathcal{S}(K)$
defined by $\sigma \circ \varphi = \varphi \circ E$

Observation

$(\mathcal{S}(K), \sigma)$ is a difference ring.

Example

- $a = 0$ in $\mathcal{S}(K)$ $\iff a_n = 0$ for all large enough n
- $a = b$ in $\mathcal{S}(K)$ $\iff a_n = b_n$ for all large enough n
- $K[n]$, $K(n)$, $\mathcal{H}(K)$ naturally embed into $\mathcal{S}(K)$

Henceforth we work in $(\mathcal{S}(K), \sigma)$.

A. Polynomial solutions

Given: $L \in K[n]\langle\sigma\rangle$, $L \neq 0$

Find: a basis of the space $\{y \in K[n]; Ly = 0\}$

Outline of algorithm

- 1 Find an upper bound for $\deg y$.
- 2 Use the method of undetermined coefficients.

B. Rational solutions

Given: $L \in K[n]\langle\sigma\rangle$, $L \neq 0$

Find: a basis of the space $\{y \in K(n); Ly = 0\}$

Outline of algorithm

- 1 Find a universal denominator for y .
- 2 Find polynomial solutions of the equation satisfied by the numerator of y .

C. Hypergeometric solutions

Given: $L = \sum_{k=0}^d p_k \sigma^k \in K[n]\langle\sigma\rangle$, $L \neq 0$

Find: a generating set for $\text{Lin}(\{y \in \mathcal{H}(K); Ly = 0\})$

Outline of algorithm

1 Construct the Riccati equation for $r = \frac{\sigma y}{y} \in K(n)$:

$$\sum_{k=0}^d p_k \prod_{j=0}^{k-1} \sigma^j r = 0 \quad (2)$$

C. Hypergeometric solutions

2 Use the ansatz

$$r = z \frac{a \sigma c}{b c}$$

with $z \in K^*$, $a, b, c \in K[n]$ **monic**,
 a, c **coprime**, $b, \sigma c$ **coprime**, $a, \sigma^k b$ **coprime** for all $k \in \mathbb{N}$:

$$\sum_{k=0}^d z^k p_k \left(\prod_{j=0}^{k-1} \sigma^j a \right) \left(\prod_{j=k}^{d-1} \sigma^j b \right) \sigma^k c = 0 \quad (3)$$

C. Hypergeometric solutions

3 Construct a finite set of candidates for (a, b, z) using the following consequences of (3):

- $a \mid p_0,$
- $b \mid \sigma^{1-d} p_d,$
- $\sum_{\substack{0 \leq k \leq d \\ \deg P_k = m}} \text{lc}(P_k) z^k = 0$

where $P_k = p_k \left(\prod_{j=0}^{k-1} \sigma^j a \right) \left(\prod_{j=k}^{d-1} \sigma^j b \right),$
 $m = \max_{0 \leq k \leq d} \deg P_k.$

C. Hypergeometric solutions

- 4 For each candidate triple (a, b, z) , find polynomial solutions c of the equation

$$\sum_{k=0}^d z^k P_k \sigma^k c = 0.$$

C. Hypergeometric solutions

Example

(AMM problem no. 10375) Solve

$$y_{n+2} - 2(2n+3)^2 y_{n+1} + 4(n+1)^2(2n+1)(2n+3)y_n = 0.$$

$$p_2(n) = 1$$

$$p_1(n) = -2(2n+3)^2$$

$$p_0(n) = 4(n+1)^2(2n+1)(2n+3)$$

C. Hypergeometric solutions

1 Riccati equation:

$$p_2(n) r(n+1)r(n) + p_1(n) r(n) + p_0(n) = 0$$

2 plug in the ansatz:

$$\begin{array}{r} z^2 \\ + z \\ + \end{array} \begin{array}{l} p_2(n) \\ p_1(n) \\ p_0(n) \end{array} \begin{array}{l} a(n+1) a(n) \\ a(n) b(n+1) \\ b(n+1) b(n) \end{array} \begin{array}{l} c(n+2) \\ c(n+1) \\ c(n) \end{array} = 0$$

C. Hypergeometric solutions

3 candidates for (a, b, z) :

- $a(n) \mid 4(n+1)^2(2n+1)(2n+3)$

- $b(n) \mid 1$

Take, e.g., $a(n) = (n+1)(n+\frac{1}{2})$ and $b(n) = 1$.

- $z^2 - 8z + 16 = (z-4)^2 = 0$

Take $z = 4$.

C. Hypergeometric solutions

4 equation for c :

$$(n+2)c(n+2) - (2n+3)c(n+1) + (n+1)c(n) = 0$$

Polynomial solution: $c(n) = 1$

$$r(n) = z \frac{a(n)}{b(n)} \frac{c(n+1)}{c(n)} = 4(n+1) \left(n + \frac{1}{2} \right)$$

$$\frac{y_{n+1}}{y_n} = r(n) = (2n+1)(2n+2) \implies y_n = C(2n)!$$