

Solvers II — Preconditioning and Deflation

Lattice Practices 2012

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Outline

Motivation

The curse of ill-conditioning

Preconditioning

Preconditioning — Basics

Preconditioned Krylov subspace methods

Preconditioners

Deflation

Summary



How to improve an optimal method?

Solvers I: Krylov subspace methods are all-duty solvers

- ▶ “Optimal” methods for any application
- ▶ Fast (i.e., short-recurrence) solvers for many applications
- ▶ Convergence dependent on conditioning of A , e.g.,
 - ▶ Conjugate Gradients

$$\|e^{(k)}\| \leq \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^k \|e^{(0)}\|_A, \quad \kappa = \frac{\lambda_{\max}(A)}{\lambda_{\min}(A)}$$

How to improve convergence of Krylov subspace methods?

1. Preconditioning
2. Deflation



Scaling issues in Numerical Simulations

Numerical simulations of **partial differential equations** (PDEs)

$$\mathcal{L}\psi = \varphi$$

Discretization of \mathcal{L} on mesh with spacing a yields

$$\mathbf{L}x = f$$

- ▶ Depending on PDE order and order of discretization

$$\kappa(\mathbf{L}) \sim a^{-\sigma}, \quad \sigma \in \mathbb{N}^+$$

- ▶ Increasing accuracy of discretization ($a \rightarrow 0$)

$$\kappa(\mathbf{L}) \longrightarrow \infty \quad (a \rightarrow 0)$$

Performance of Krylov methods **deteriorates** when $a \rightarrow 0$!



Preconditioning — Idea

Idea: Improve conditioning of A in $Ax = b$!

- ▶ Instead of solving $Ax = b$ consider solving

$$\begin{aligned}S_\ell A S_r y &= S_\ell b \\ x &= S_r y\end{aligned}$$

with **preconditioners** S_ℓ, S_r s.t. $\kappa(S_\ell A S_r) \ll \kappa(A)$

Open questions

- ▶ What are the design goals for preconditioners?
- ▶ What are suitable choices of S_ℓ, S_r ?
- ▶ How does the preconditioner fit in the iteration
 - ▶ Ideally only $A\cdot, S_\ell\cdot$ and $S_r\cdot$ are required

For now consider only left-preconditioning with $S = S_\ell$



Preconditioning — Observations

Considering extreme cases

- ▶ $S = I$
⇒ $SA = A$ original setting
- ▶ $S = A^{-1}$
⇒ $SA = I$ and $\kappa(SA) = 1$ (ideal)
- ▶ $S = A^\dagger$
⇒ $SA = A^\dagger A$ hermitian, but $\kappa(SA) = \kappa(A)^2$

In order to speed up convergence preconditioner S should

- ▶ approximate A^{-1}
- ▶ be cheap to compute ($S \cdot$)





Preconditioning — CG

Recall: Conjugate Gradients requires A hermitian

Problem: SA in general no longer hpd even if S is hpd, but

$$\langle SAx, y \rangle_{S^{-1}} = \langle Ax, y \rangle_2 = \langle x, Ay \rangle_2 = \langle x, SAy \rangle_{S^{-1}}$$

Solution: Replace all $\langle \cdot, \cdot \rangle_2$ by $\langle \cdot, \cdot \rangle_{S^{-1}}$

- ▶ Rewriting the algorithm you can get rid of $\langle \cdot, \cdot \rangle_{S^{-1}}$ again
- ▶ CG variants exist for any A hermitian in some $\langle \cdot, \cdot \rangle_B$

Similar arguments can be made for all methods with prerequisites, e.g., **MINRES**, SUMR





PCG — Algorithm

Preconditioned Conjugate Gradients

$$r^{(0)} = b - Ax^{(0)}, z^{(0)} = Sr^{(0)}, p^{(0)} = z^{(0)}$$

for $k = 1, 2, \dots$ **do**

$$\alpha_{k-1} = \frac{\langle r^{(k-1)}, z^{(k-1)} \rangle_2}{\langle Ap^{(k-1)}, p^{(k-1)} \rangle_2}$$

$$x^{(k)} = x^{(k-1)} + \alpha_{k-1} p^{(k-1)}$$

$$r^{(k)} = r^{(k-1)} - \alpha_{k-1} Ap^{(k-1)}$$

$$z^{(k)} = Sr^{(k)}$$

$$\beta_{k-1} = \frac{\langle r^{(k)}, z^{(k)} \rangle_2}{\langle r^{(k-1)}, z^{(k-1)} \rangle_2}$$

$$p^{(k)} = z^{(k)} + \beta_{k-1} p^{(k-1)}$$

end for





Preconditioned GMRES(m)

$$r^{(0)} = S(b - Ax^{(0)}), \beta = \|r^{(0)}\|_2, v_1 = \beta^{-1}r^{(0)}$$

while not converged **do**

for $j = 1, \dots, m$ **do**

$$w = SAV_j$$

for $i = 1, \dots, j$ **do**

$$h_{i,j} = \langle w, v_j \rangle_2$$

$$w = w - h_{i,j}v_j$$

end for

$$h_{j+1,j} = \|w\|_2$$

$$v_{j+1} = h_{j+1,j}^{-1}w$$

end for

 Define $V_m = [v_1 \mid \dots \mid v_m]$, $H_{m+1,m} = \{h_{i,j}\}_{1 \leq j \leq m, 1 \leq i \leq j+1}$

 Solve $y_m = \operatorname{argmin}_y \|\beta e_1 - H_{m+1,m}y\|_2$

$$x^{(0)} = x^{(0)} + V_m y_m$$

end while





Preconditioned BiCGstab

$$r^{(0)} = b, \beta_0 = 0$$

$$\hat{r} = r$$

shadow residual $\langle p, \hat{p} \rangle_2 \neq 0$

for $k = 0, 1, \dots$ **do**

$$\rho_k = \langle r^{(k)}, \hat{r} \rangle_2$$

$$\beta_k = \frac{\rho_k}{\rho_{k-1}} \cdot \frac{\alpha_{k-1}}{\omega_{k-1}}$$

$$p^{(k)} = r^{(k)} + \beta_k(p^{k-1} - \omega_{k-1}v^{(k-1)})$$

$$\hat{p}^{(k)} = Sp^{(k)}$$

$$\alpha_k = \frac{\rho_k}{\langle A\hat{p}^{(k)}, \hat{r} \rangle_2}$$

$$x^{(k+\frac{1}{2})} = x^{(k)} + \alpha_k \hat{p}^{(k)}$$

$$s^{(k)} = r^{(k)} - \alpha_k A\hat{p}^{(k)}$$

$$s^{(k)} \equiv r^{(k+\frac{1}{2})}$$

$$\hat{s}^{(k)} = Ss^{(k)}$$

$$\omega_k = \frac{\langle s^{(k)}, A\hat{s}^{(k)} \rangle_2}{\langle A\hat{s}^{(k)}, A\hat{s}^{(k)} \rangle_2}$$

$$x^{(k+1)} = x^{(k+\frac{1}{2})} + \omega_k \hat{s}^{(k)}$$

$$r^{(k+1)} = s^{(k)} - \omega_k A\hat{s}^{(k)}$$

end for



Preconditioners

Aims for the construction of **preconditioners** S

1. $S \approx A^{-1}$ to get speed-up
2. $S \cdot$ should be cheap (1 application per iterate)

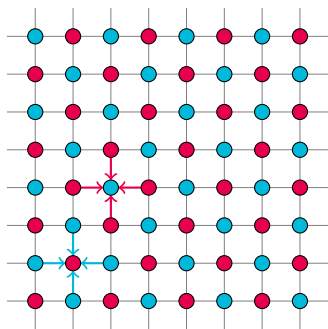
Classes of preconditioners to be discussed

- ▶ Structural preconditioners
- ▶ Splitting-based preconditioners
- ▶ Domain decomposition preconditioners
- ▶ Multigrid preconditioners
- ▶ Incomplete decomposition preconditioners



Odd-even preconditioning

Discretizations on lattices with next neighbor coupling



- ▶ Nodes are **odd** or **even**

Ordering by **odd-even**

$$A = \begin{bmatrix} A_{oo} & A_{oe} \\ A_{eo} & A_{ee} \end{bmatrix}$$

with diagonal A_{oo} and A_{ee}

- ▶ A_{oo}^{-1}, A_{ee}^{-1} trivial
- ▶ **odd** decoupled
- ▶ **even** decoupled

Solve first **even** then **odd**



Odd-even preconditioning

With $\hat{A}_{ee} = A_{ee} - A_{eo}A_{oo}^{-1}A_{oe}$ solution of $Ax = b$ given by

Odd-Even Reduction

$$y_o = A_{oo}^{-1}b_o$$

$$\text{Solve } \hat{A}_{ee}x_e = b_e - A_{eo}y_o$$

$$x_o = y_o - A_{oo}^{-1}A_{oe}x_e$$

- ▶ Approximately solving $\hat{A}_{ee}x_e = b_e - A_{eo}y_o$
⇒ Odd-Even preconditioner
- ▶ If A has constant diagonal $\kappa(\hat{A}_{ee}) < \kappa(A)$
⇒ Solving \hat{A}_{ee} is easier than solving A
- ▶ Since A_{oo}^{-1} is cheap (diagonal!)
⇒ Cost for $\hat{A}_{ee} \cdot \approx$ Cost for A .



Splitting methods

Splitting methods use the **additive** decomposition of A

- ▶ Jacobi iteration $x^{(k+1)} = x^{(k)} + D^{-1}r^{(k)}$
- ▶ Gauss-Seidel iteration $x^{(k+1)} = x^{(k)} + (D + L)^{-1}r^{(k)}$
- ▶ Successive over-relaxation

$$x^{(k+1)} = (1 - \omega)x^{(k)} + \omega(D + L)^{-1}r^{(k)}$$

- ▶ In general if $A = M + N$

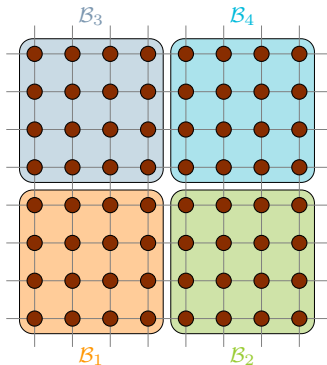
$$x^{(k+1)} = x^{(k)} + M^{-1}r^{(k)} = x^{(k)} + M^{-1}Ae^{(k)}$$

Convergent iff $\|I - M^{-1}A\| < 1$



Domain Decomposition*

- ▶ Split the computational domain into subdomains \mathcal{B}_i
- ▶ Solve system iteratively on each subdomain



- ▶ Canonical injection \mathcal{I}_j

$$\mathcal{I}_j e_i = e_{(\mathcal{B}_j)_i}$$

- ▶ Restriction of x onto \mathcal{B}_j

$$x_{\mathcal{B}_j} = \mathcal{I}_j^\dagger x$$

- ▶ Restriction of A onto \mathcal{B}_j

$$A_{\mathcal{B}_j} = \mathcal{I}_j^\dagger A \mathcal{I}_j$$

*Domain decomposition dates back to H. Schwarz (1870)



Additive and Multiplicative Schwarz

Additive Schwarz

```
for  $k = 0, 1, \dots$  do  
   $r^{(k)} = b - Ax^{(k)}$   
  for  $j = 1, 2, \dots, n_B$  do  
     $x_{\mathcal{B}_j}^{(k+1)} = x_{\mathcal{B}_j}^{(k)} + A_{\mathcal{B}_j}^{-1} r_{\mathcal{B}_j}^{(k)}$   
  end for  
end for
```

- ▶ Block-Jacobi
- ▶ Embarrassingly Parallel

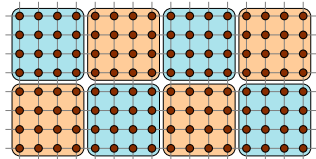
Schwarz methods in general

- ⊕ Data parallel
- ⊕ Computation parallel

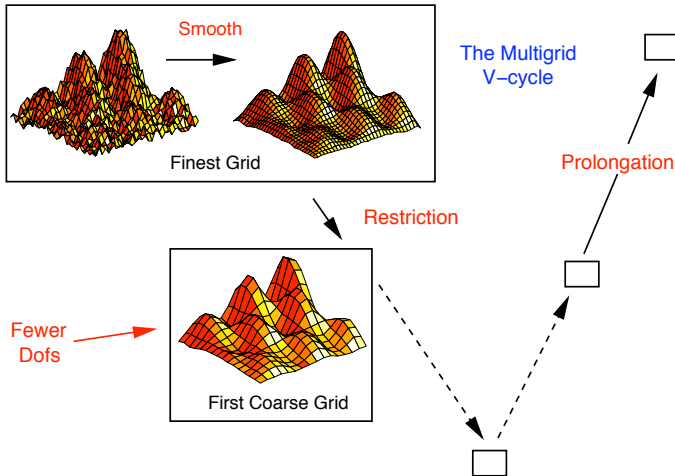
Multiplicative Schwarz

```
for  $k = 0, 1, \dots$  do  
  for  $j = 1, 2, \dots, n_B$  do  
     $r = b - Ax$   
     $x_{\mathcal{B}_j} = x_{\mathcal{B}_j} + A_{\mathcal{B}_j}^{-1} r_{\mathcal{B}_j}$   
  end for  
end for
```

- ▶ Block-Gauss-Seidel
- ▶ Sequential (\rightarrow coloring)



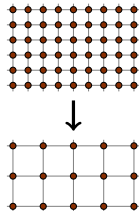
Multigrid



(Algebraic) Multigrid

- Given:**
- ▶ $Ax = b$
 - ▶ Iterative method S

- Wanted:**
- ▶ Hierarchy of systems
 $A_\ell x_\ell = b_\ell, \quad \ell = 0, \dots, L$
 - ▶ Intergrid transfer operators
 $P_{\ell+1}^\ell : \mathbb{C}^{n_{\ell+1}} \rightarrow \mathbb{C}^{n_\ell}$



Smoother

$$S_\ell : \mathbb{C}^{n_\ell} \rightarrow \mathbb{C}^{n_\ell}$$

“High modes”

Interpolation

$$P_{\ell+1}^\ell : \mathbb{C}^{n_{\ell+1}} \rightarrow \mathbb{C}^{n_\ell}$$

“Low modes”

Complementarity of Smoother and Interpolation



Generic Multigrid Algorithm — $MG_\ell(A_\ell, b_\ell)$

if $\ell = L$ **then**

$$x_L = A_L^{-1}b_L$$

else

$$x_\ell = 0$$

for $i = 1, \dots, \nu_1$ **do**

$$x_\ell = S(x_\ell, b_\ell)$$

“Pre-smoothing”

end for

$$x_{\ell+1} = MG(A_{\ell+1}, R_{\ell+1}^\ell(b_\ell - Ax_\ell))$$

$$x_\ell = x_\ell + P_{\ell+1}^\ell x_{\ell+1}$$

“Coarse-grid correction”

for $i = 1, \dots, \nu_2$ **do**

$$x_\ell = S(x_\ell, b_\ell)$$

“Post-smoothing”

end for

end if



Incomplete LU

Recall: Direct methods are based on factorization of A

$$A = L \cdot U$$

Drawback: Fill-In in L and U for sparse A

Idea: Incomplete factorizations with sparse L and U

1. Prescribe the non-zero pattern (e.g., non-zeroes of A)
 - ▶ Minimize the error-matrix E in $A = \tilde{L}\tilde{U} + E$
2. Use drop-tolerance θ to drop small entries in L and U
 - ▶ If A is sparse: $(A^{-1})_{i,j} \sim \alpha^{\text{dist}_G(i,j)}$, $\alpha < 1$
 - ⇒ If i is “far” from j , L_{ij} or U_{ij} are likely dropped

ILU is a **black-box** preconditioner



Flexible Krylov subspace methods

In many cases the preconditioner is again an iterative process

- ▶ S changes in each iteration $\rightarrow S = S_k$
- ▶ There is no longer a Krylov subspace defined by

$$\mathcal{K}_k(SA, b) = \{b, SAb, (SA)^2b, \dots, (SA)^{k-1}b\}$$

\Rightarrow Convergence theory does not hold anymore

- ▶ Krylov subspace methods have to be modified!
 \Rightarrow Flexible Krylov subspace methods



Flexible CG — Algorithm

Flexible Conjugate Gradients

$$r^{(0)} = b - Ax^{(0)}, z^{(0)} = S_0 r^{(0)}, p^{(0)} = z^{(0)}$$

for $k = 1, 2, \dots$ **do**

$$\alpha_{k-1} = \frac{\langle r^{(k-1)}, z^{(k-1)} \rangle_2}{\langle Ap^{(k-1)}, p^{(k-1)} \rangle_2}$$

$$x^{(k)} = x^{(k-1)} + \alpha_{k-1} p^{(k-1)}$$

$$r^{(k)} = r^{(k-1)} - \alpha_{k-1} Ap^{(k-1)}$$

$$z^{(k)} = S_k r^{(k)}$$

$$\beta_{k-1} = \frac{\langle r^{(k)} - r^{(k-1)}, z^{(k)} \rangle_2}{\langle r^{(k-1)}, z^{(k-1)} \rangle_2}$$

$$p^{(k)} = z^{(k)} + \beta_{k-1} p^{(k-1)}$$

end for

- ▶ If $S_k = S$ for all $k \implies z^{(k)} \perp r^{(k-1)}$
- ▶ Flexible CG preserves local optimality



Flexible GMRES(m)

$$r^{(0)} = S_0(b - Ax^{(0)}), \beta = \|r^{(0)}\|_2, v_1 = \beta^{-1}r^{(0)}$$

while not converged **do**

for $j = 1, \dots, m$ **do**

$$z_j = S_j v_j$$

$$w = Az_j$$

for $i = 1, \dots, j$ **do**

$$h_{i,j} = \langle w, v_j \rangle_2$$

$$w = w - h_{i,j} v_j$$

end for

$$h_{j+1,j} = \|w\|_2$$

$$v_{j+1} = h_{j+1,j}^{-1} w$$

end for

 Define $Z_m = [z_1 \mid \dots \mid z_m]$, $H_{m+1,m} = \{h_{i,j}\}_{1 \leq j \leq m, 1 \leq i \leq j+1}$

 Solve $y_m = \operatorname{argmin}_y \|\beta e_1 - H_{m+1,m} y\|_2$

$$x^{(0)} = x^{(0)} + Z_m y_m$$

end while





Preconditioners — Summary

Preconditioning improves convergence if $\kappa(SA) \ll \kappa(A)$

- ▶ There is a wide variety of preconditioners available
 - ▶ Most of them require knowledge about A or its origins
- ▶ Aims for the construction of preconditioners S are
 - ▶ $S \approx A^{-1}$ and $S \cdot$ cheap

Preconditioning makes Krylov subspace methods more robust

- ▶ Reducing $\kappa(A)$ helps controlling the error $e^{(k)}$, since

$$\|e\|_2 \leq c\kappa(A)^{-1}\|r\|_2$$

⇒ If $\kappa(A) \gg 1$ results based on $\|r\|_2$ should not be trusted!

⇒ If $\kappa(A) \gg 1$ a preconditioner is **mandatory!**



Deflation — Idea

Typically convergence is slowed down by **small eigenmodes**

- ▶ Given the “troublesome” modes v_1, \dots, v_ℓ
 \Rightarrow **deflate** the subspace $\mathcal{V} = \text{colspan}(\underbrace{[v_1 \mid \dots \mid v_\ell]}_{=V})$

Similar to preconditioning, instead of $Ax = b$ solve

$$\begin{aligned}(I - \pi_A(V)) A \hat{x} &= (I - \pi_A(V)) b \\ x &= \hat{x} + \pi_A(V) b\end{aligned}$$

with $\pi_A(V) = V(V^\dagger AV)^{-1}V^\dagger A$

- ▶ In case v_i are eigenmodes, $V^\dagger AV = \text{diag}(\lambda_1, \dots, \lambda_\ell)$
 $\Rightarrow (V^\dagger AV)^{-1}$ nothing to worry about



Deflation — Conjugate Gradients Theory

The **effective condition number** κ_{eff} replaces κ in theory

$$\begin{aligned}\kappa_{\text{eff}} &= \frac{\mu_1}{\mu_\ell} \\ \mu_1 &= \max_{x \neq 0} \frac{\langle A(I - \pi_A(V))x, x \rangle_2}{\langle x, x \rangle_2} \\ \mu_\ell &= \min_{x \in \mathcal{V}^\perp \setminus \{0\}} \frac{\langle A(I - \pi_A(V))x, x \rangle_2}{\langle x, x \rangle_2}\end{aligned}$$

- ▶ If v_i are smallest ℓ eigenmodes

$$\kappa_{\text{eff}} = \frac{\lambda_{\max}}{\lambda_{\ell+1}}$$

where $\lambda_{\ell+1}$ is the $(\ell + 1)^{\text{st}}$ smallest EW



Deflated CG — Algorithm

Deflated CG (Deflation space $\mathcal{V} = \text{colspan}(V)$)

$$x^{(0)} = x^{(0)} + \pi_A(V)b$$

$$r^{(0)} = b - Ax^{(0)}$$

$$p^{(0)} = (I - \pi_A(V))r^{(0)}$$

for $k = 1, 2, \dots$ **do**

$$\alpha_{k-1} = \frac{\langle r^{(k-1)}, r^{(k-1)} \rangle_2}{\langle Ap^{(k-1)}, p^{(k-1)} \rangle_2}$$

$$x^{(k)} = x^{(k-1)} + \alpha_{k-1}p^{(k-1)}$$

$$r^{(k)} = r^{(k-1)} - \alpha_{k-1}Ap^{(k-1)}$$

$$\beta_{k-1} = \frac{\langle r^{(k)}, r^{(k)} \rangle_2}{\langle r^{(k-1)}, r^{(k-1)} \rangle_2}$$

$$p^{(k)} = (I - \pi_A(V))r^{(k)} + \beta_{k-1}p^{(k-1)}$$

end for



GMRES(m)

On restart all information about $\mathcal{K}_m(A, r^{(0)})$ is lost!

- ▶ Use deflation technique to transfer information

Note: Due to the **Arnoldi relation** $V_m^\dagger AV_m = H_{m,m}$

$$\min_{x \in \mathcal{K}_m(A, r^{(0)}), x \neq 0} \frac{\langle Ax, x \rangle_2}{\langle x, x \rangle_2} = \min_{y \neq 0} \frac{\langle H_{m,m} y, y \rangle_2}{\langle y, y \rangle_2}$$

- ▶ Eigenmodes w_1, \dots, w_m of $H_{m,m}$ give approximations

$$V_m^\dagger (AV_m w_i - \lambda_i V_m w_i) = 0$$

- ▶ Vectors $V_m w_i$ are called **Ritz vectors** (\rightarrow ARPACK)

Idea: Use smallest eigenmodes of $H_{m,m}$ in deflation



Deflated GMRES(m) — Sketch**for** $\ell = 0, 1, \dots$ **do**

$$r^{(0)} = b - Ax^{(0)}, \beta = \|r^{(0)}\|_2, v_1 = \beta^{-1}r^{(0)}, \tilde{V} = \emptyset$$

Compute $V_m, H_{m+1,m}$ based on initial \tilde{V} (Arnoldi)Compute smallest Ritz vectors $V_m w_1, \dots, V_m w_\ell$

$$y_m = \operatorname{argmin}_y \|\beta e_1 - H_{m+1,m} y\|_2$$

$$x^{(0)} = x^{(0)} + V_m y_m$$

$$\tilde{V} = [V_m w_1 \mid \dots \mid V_m w_\ell]$$

end for

- ▶ For a more detailed description see [3]
- ▶ Reusing information upon restart is also known as...
 - ▶ ...recycling
 - ▶ ...augmenting



Deflation — Summary

Deflation “hides” most difficult part of the problem

- ▶ Computation of eigenmodes necessary
 - ▶ possibly on-the-fly (Deflated GMRES(m))
 - ▶ possibly a priori knowledge available
 - ▶ approximations viable (\rightarrow ARPACK)
- ▶ Analysis of general deflation subspaces \mathcal{V} (cf. [2])

Eigenmode **deflation** suffers from scaling (i.e., $a \rightarrow 0$)

- ▶ In order to have constant number of iterations for $a \rightarrow 0$

$$\kappa_{\text{eff}} = \text{const} \iff \lambda_{\min}^{\text{eff}} > \sigma$$

- ▶ Often number N_σ of EW below threshold σ fulfills

$$N_\sigma \sim V \rightarrow \infty \quad (a \rightarrow 0)$$

\Rightarrow More eigenmodes need to be computed as $a \rightarrow 0$



Summary

To find an efficient solver is hard, but there are **guidelines**

- ▶ Use as much information about your system as possible
 - ▶ In the choice of the Krylov subspace method
 - ▶ Short recurrence method available?
 - ▶ Optimal method available?
 - ▶ In the choice of the preconditioner
- ▶ Adjust parameters of your method w.r.t. hardware, e.g.,
 - ▶ Restart length in GMRES(m)
 - ▶ Dimension of the deflation subspace
 - ▶ Dimension of the subdomains in domain decomposition

Most often there is no obvious optimal choice for the solver!

Construction of optimal solvers is ongoing research!





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